

**$K_4 - e$  Designs with a Hole**

by

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## Abstract

In this paper we look at  $K_4 - e$  designs on  $K_{w-v} + v$ . We settle the case when  $w$  and  $v$  are of the same parity.

## Acknowledgments

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## Chapter 1

### Introduction

A  $G$ -design of  $H$  is an edge-disjoint decomposition of  $H$  into isomorphic copies of the graph  $G$ . When  $H$  is the complete graph on  $n$  vertices,  $H = K_n$ , it is called a  $G$ -design of order  $n$ . The *spectrum problem* for  $G$ , i.e. solving for which values of  $n$  there exists a  $G$ -design of order  $n$ , has been solved for many graphs  $G$ , including all graphs  $G$  on less than 6 vertices. [1, 2] A *complete graph with a hole of size  $v$* ,  $K_n \setminus K_v$ , is a complete graph on  $n$  vertices where the edges of a complete subgraph of size  $v$  have been removed.

Let  $V$  be an independent set of vertices of size  $v$ . The graph  $G + v$  is a graph  $G$  where every vertex in  $V$  is adjacent to every vertex in  $G$ . So,  $K_d + v = K_n \setminus K_v$  where  $d = n - v$ .

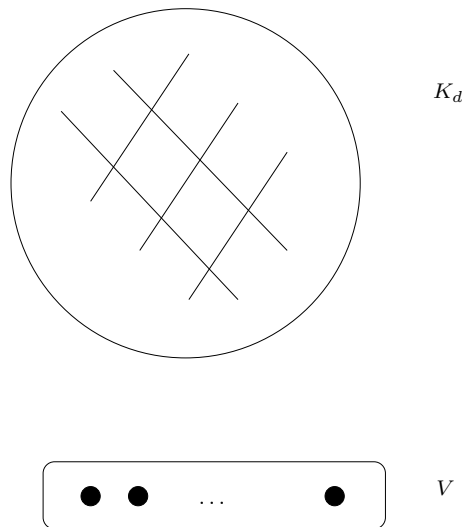


Figure 1.1:  $K_d + v$

$G$ -designs with holes of size  $v$  partition the edges of  $K_d + v$  into edge disjoint copies of  $G$ . These designs were first studied by Doyen and Wilson in 1973 where  $G = K_3$ . [3]

A *Steiner Triple System* of order  $v$ ,  $STS(v)$ , is an ordered pair  $(S, T)$ , where  $S$  is a finite set of points of size  $v$ , and  $T$  is a set of 3-element subsets of  $S$  called triples, where any two distinct points lie in exactly one triple.  $(S_1, T_1)$  is a *subsystem* of  $(S, T)$  if,  $S_1 \subseteq S$ ,  $T_1 \subseteq T$ , and  $(S_1, T_1)$  is a  $STS$ .

Doyen and Wilson's theorem give the necessary and sufficient conditions of a  $STS(w)$  which contains a  $STS(v)$  as a subsystem to be  $w \equiv 1$  or  $3 \pmod{6}$  and  $v \equiv 1$  or  $3 \pmod{6}$  and  $w \geq 2v + 1$ . Since the structure of the subsystem plays no role in the design on  $w$ , the subsystem could just be a hole. Thus we have a  $K_3$  design on  $K_{w-v} + v$ . These also exist when  $w \equiv v \equiv 5 \pmod{6}$ .

This result has been extended to find  $G$ -designs on graphs with holes where  $G$  is a cycle of length at most 14 and a triangle with a pendent edge. [4, 5]

A *packing* of  $K_n$  with a graph  $G$  is a triple,  $(S, T, L)$ , where  $S$  is the vertex set of  $K_n$ ,  $T$  is a collection of edge disjoint copies of  $G$  from the edge set of  $K_n$ , and  $L$  is the collection of edges in  $K_n$  not belonging to one of the copies of  $G$  in  $T$ . The collection of edges in  $L$  is called the *leave*. When  $|L|$  is as small as possible,  $(S, T, L)$  is called a *maximum packing of order  $n$* .

We look at the case when  $G = K_4 - e$ , the complete graph on four vertices minus one edge. We will denote a  $K_4 - e$ , as shown in Figure 1.2, by any of the quadruples  $(a, b, c, d)$ ,  $(a, b, d, c)$ ,  $(b, a, c, d)$  or  $(b, a, d, c)$ .

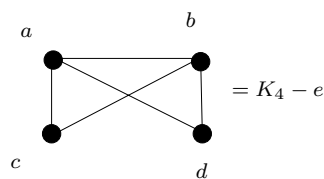


Figure 1.2:  $K_4 - e$

In 1977, Bermond and Schonheim showed a  $K_4 - e$  design of order  $n$  exists if and only if  $n \equiv 0$  or  $1 \pmod{5}$  and  $n \geq 6$ . Hoffman, Lindner, Sharry, and Street, in 1993, solved

maximum packings of  $K_n$ , with  $K_4 - e$ . [6] These two results solve the  $K_4 - e$  design on  $K_d + v$  when  $v = 0, 1, 2, 3$ . (A  $K_4 - e$  design of order  $n$  can be considered a  $K_n + 0$  or a  $K_{n-1} + 1$ .)

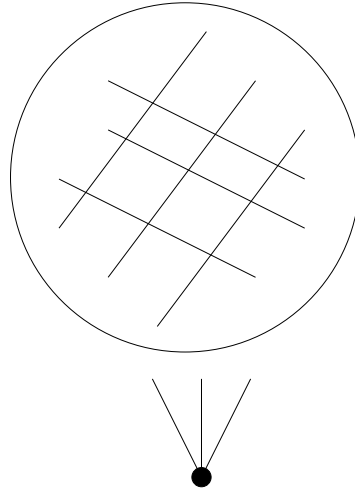


Figure 1.3:  $K_{n-1} + 1$

The maximum packing paper tells us precisely when  $K_n$  has a leave of a single edge or a triangle. These graphs may be considered a  $K_{n-2} + 2$  and  $K_{n-3} + 3$  respectively.

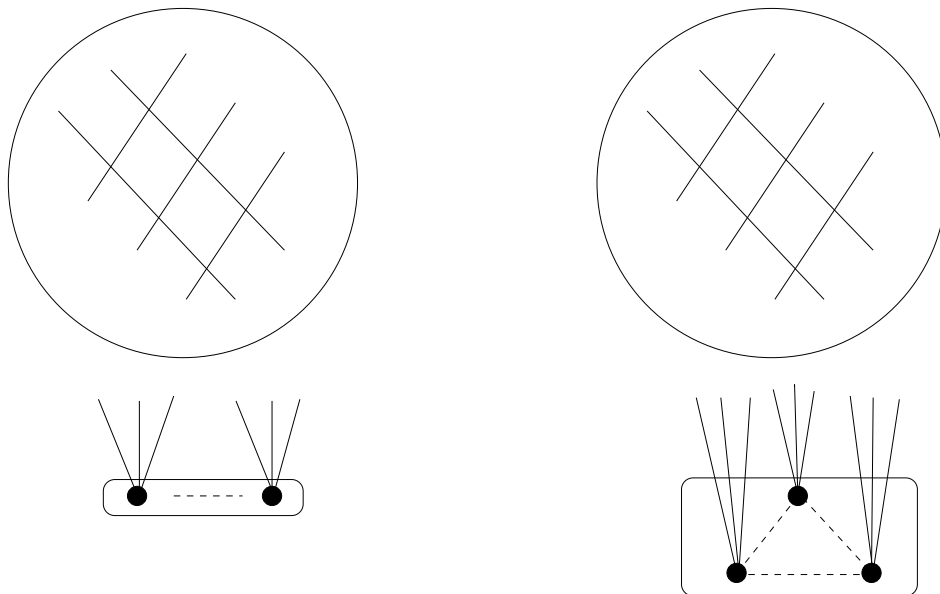


Figure 1.4:  $K_{n-2} + 2$  and  $K_{n-3} + 3$



The two main tools we use are difference methods and 1-factors. When  $d$  is even let  $d = 2t$ . We will depict  $K_d$  as  $\mathbb{Z}_t \times \{1, 2\}$  with all possible edges between the sets (see Figure 1.5). We refer to the edges and vertices in  $K_d$  as being upstairs and the vertices in  $V$  as downstairs.

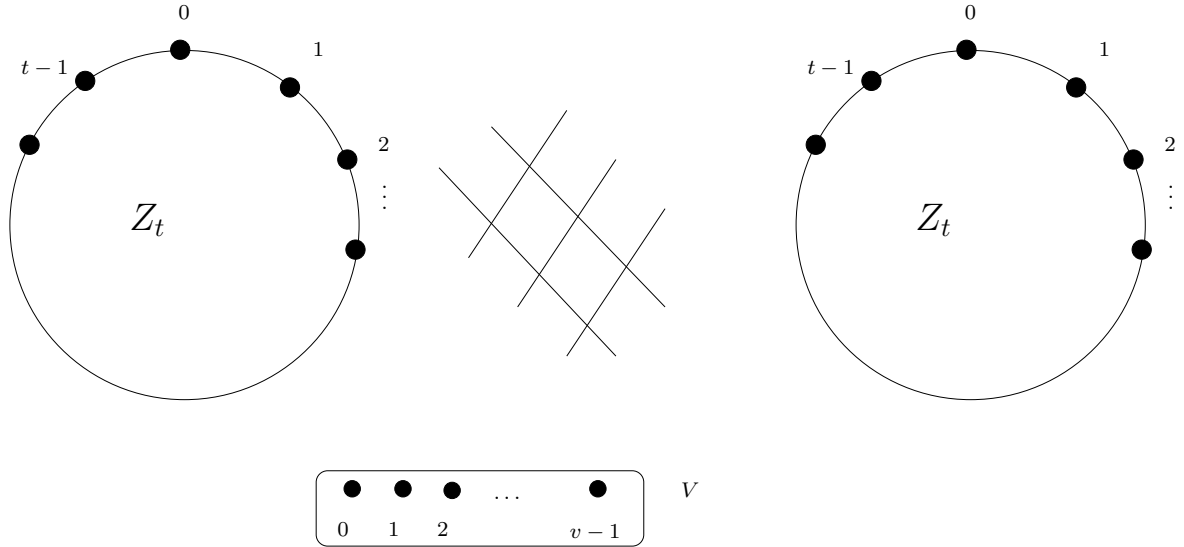


Figure 1.5:  $\mathbb{Z}_t \times \{1, 2\} + v$

If  $a$  and  $b$  are integers, define  $|b - a|_t$ , the *difference of  $a$  and  $b$  (mod  $t$ )*, to be the smallest non-negative integer congruent to  $b - a$  or  $a - b$  (mod  $t$ ). So,  $0 \leq |b - a|_t \leq t/2$ . The edges of the complete graph on vertex set  $\mathbb{Z}_t$  are partitioned into differences  $1, 2, \dots, \lfloor t/2 \rfloor$  and the edge  $ab$  is defined to be of *pure difference*  $|b - a|_t$ . When  $t/2$  is an integer this is called the *half difference* and will generate  $t/2$  edges in each  $\mathbb{Z}_t$ . All other pure differences will generate  $t$  edges in each  $\mathbb{Z}_t$ . *Mixed differences* describe the distance between vertices in different sets. Let  $(x, 1)$  be a vertex in  $\mathbb{Z}_t \times \{1\}$  and let  $(y, 2)$  be a vertex in  $\mathbb{Z}_t \times \{2\}$ . The mixed difference between them is  $y - x$  (mod  $t$ ). There are  $t$  mixed differences each consisting of  $t$  edges.

Let  $D_t = \{1, 2, \dots, \lfloor t/2 \rfloor\}$ , the set of differences in  $\mathbb{Z}_t$ . If  $S$  is a subset of  $D_t$ , we define  $G[S]$  to be the graph on  $t$  vertices induced by the differences in  $S$ . We call  $x \in D_t$  a *good difference* if  $t/\gcd(x, t)$  is even.

A *1-factor* is a perfect matching, i.e. a set of a single edges that contain each vertex exactly once. When the edges of a graph can be partitioned into a set of 1-factors, this is called a *1-factorization*.

**Lemma 1.1 (Stern and Lenz)** *Let  $\emptyset \neq S \subseteq D_t$ .  $G[S]$  has a 1-factorization if and only if  $S$  contains at least one good difference.*

In the proof they use:

**Lemma 1.2** *Let  $G$  be a simple regular graph and  $G'$  an isomorphic copy of  $G$ . Form the graph  $H$  by adding an edge between each vertex in  $G$  and its isomorphic mate in  $G'$ . Then  $H$ , the graph shown below, has a 1-factorization.*

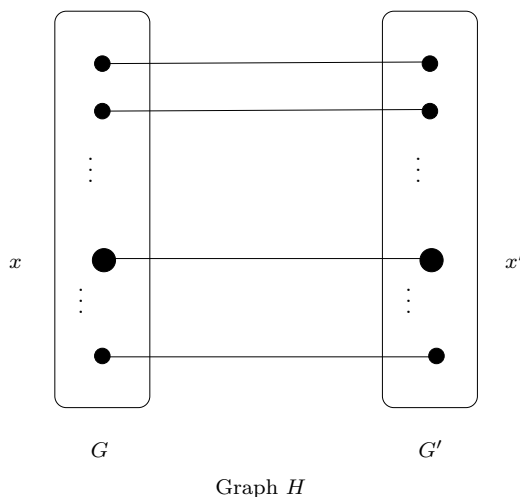


Figure 1.6: Graph  $H$

To ensure the use of these lemmas is possible while constructing blocks, we will use the same pure differences in each  $\mathbb{Z}_t \times \{1\}$  and  $\mathbb{Z}_t \times \{2\}$  and make sure we have at least one mixed difference remaining.

Let  $S, U \subseteq D_t$ , let  $T \subseteq \mathbb{Z}_d$ . Define the graph  $G[S, T, U]$  on vertex set  $\mathbb{Z}_t \times \{1, 2\}$  as follows: It's edges are those edges on  $\mathbb{Z}_t \times \{1\}$  induced by the pure differences in  $S$ , edges

on  $\mathbb{Z}_t \times \{2\}$  are induced by the pure differences in  $U$ , and those edges between  $\mathbb{Z}_t \times \{1\}$  and  $\mathbb{Z}_t \times \{2\}$  induced by the mixed differences in  $T$ . Since the edges of any single mixed difference are a 1-factor, Lemma 1.2 implies that  $G[S, T, S]$  has a 1-factorization whenever  $T \neq \{\emptyset\}$ .

In a  $K_4 - e$  design on  $K_d + v$ , there are four possible types of blocks,  $\alpha, \beta, \gamma$ , and  $\delta$ .

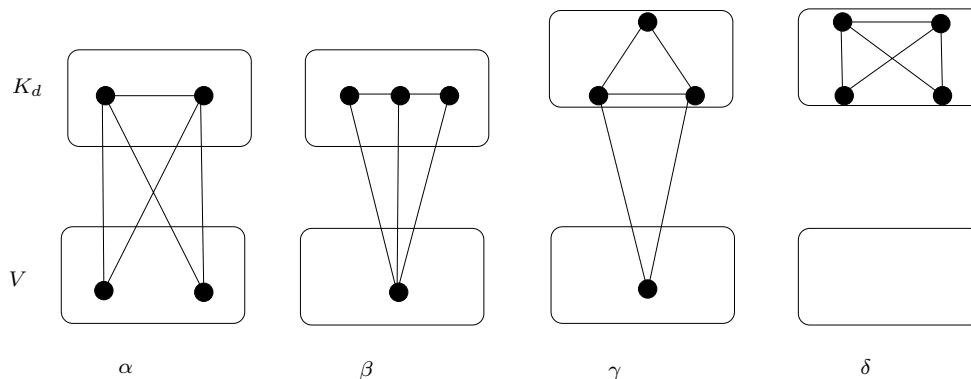


Figure 1.7: Block types

Counting the edges of each block type gives us two helpful equations. In  $K_d$ , each  $\alpha$  block contributes one edge,  $\beta$  blocks 2 edges,  $\gamma$  blocks 3 edges, and  $\delta$  blocks 5 edges. Let  $A, B, \Gamma$ , and  $\Delta$  represent the number of blocks of each respective type. Thus  $A + 2B + 3\Gamma + 5\Delta = (1/2)d(d-1)$ . Counting the edges contributed by each block type between sets  $v$  and  $d$  yields  $4A + 3B + 2\Gamma = vd$ . We also know that  $A, B, \Gamma, \Delta \geq 0$ .

In every construction we use Lemma 1.2 to find 1-factors and create  $\alpha$  blocks. Each 1-factor of  $K_d$  can be used to create  $t$   $\alpha$ 's that will contain all the edges between  $K_d$  and two vertices in  $V$ . Suppose  $a$  and  $b$  are vertices in  $V$ . For each  $xy$  in a particular 1-factor, we create the  $\alpha$  block  $(x, y, a, b)$ .

We will construct many blocks using difference methods (mod  $t$ ). In the constructions, all  $\beta$  and  $\gamma$  blocks will be defined explicitly. Blocks of type  $\alpha$  and  $\delta$  will most often be described as *base blocks*. Base blocks are  $K_4 - e$  quadruples, where initial values for the vertices in  $\mathbb{Z}_t$  will be given. Base blocks represent a set of  $t$  total blocks. They are developed

(mod  $t$ ) by taking each vertex that is an element of  $\mathbb{Z}_t$  and adding the integers (mod  $t$ ). Let  $(a, b, c, d)$  be an  $\alpha$  base block where,  $a, b \in \mathbb{Z}_t$  and  $c$  and  $d$  are vertices in  $v$ . The  $\alpha$  base block is developed (mod  $t$ ) by taking all blocks  $(a + i, b + i, c, d)$  where  $0 \leq i \leq t - 1$ . Let  $(x, y, w, z)$  be a  $\delta$  base block where  $x, y, w, z \in \mathbb{Z}_t$ . The  $\delta$  base block is developed (mod  $t$ ) by taking all blocks  $(x + i, y + i, w + i, z + i)$  where  $0 \leq i \leq t - 1$ .

Let's consider some necessary conditions for a  $K_4 - e$  design on  $K_d + v$ . First, the number of edges  $K_4 - e$  must divide the number of edges of  $K_d + v$ .

$$5|d(d + 2v - 1). \quad (1.1)$$

Second,  $v$  cannot be too large. Each block must use at least one edge upstairs. The number of blocks must be less than or equal to the number of edges in  $K_d$ . This simplifies to

$$v \leq 2(d - 1). \quad (1.2)$$

When  $d$  is odd, each vertex in  $V$  will be of odd degree. The only block type that uses an odd number of edges incident to a vertex in  $V$  is  $\beta$ . Each  $\beta$  block uses 2 edges upstairs and each vertex in  $V$  must be contained in at least one  $\beta$  block. This gives us another necessary condition when  $d$  is odd:

$$v \leq \frac{2d(d - 1)}{d + 5}. \quad (1.3)$$

Although  $K_5 + 2$  satisfies the above three necessary conditions, it was proved in the maximum packing paper that the design does not exist. Therefore, when  $d$  is odd, these conditions are not sufficient.

When  $d$  is even, 1-factors do not exist. Therefore, many of our constructions do not apply when  $d$  is odd. Thus, the main purpose (of this dissertation) is to settle the case when  $d$  is even, i.e.  $d = 2t$ . This is our main theorem:

**Theorem 1.1:** *There exists a  $K_4 - e$  design on  $K_d + v$  when  $d$  is even if and only if*

1.  $5|d(d + 2v - 1)$
2.  $v \leq 2(d - 1)$  and  $v \neq 2d - 3, v \neq 2d - 4$

## Chapter 2

### Recursion

Let  $W = \{(d, v) : \text{there exists a } K_4 - e \text{ design on } K_d + v\}$ . Recall we are now assuming  $d = 2t$ .

Part of our proof of the Main Theorem will be by induction on  $d + v$ , using the following Lemma and Theorem:

**Lemma 2.1** *If  $(d - x, v + x)$  and  $(x, v) \in W$  then  $(d, v) \in W$ .*

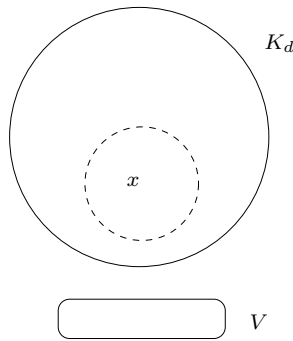


Figure 2.1: Recursion

**Proof** The union of the blocks from the designs on  $K_{d-x} + (v + x)$  and  $K_x + v$  form a design on  $K_d + v$ . □

Let  $R = \{(d, v) : 5|d(d + 2v - 1), d \text{ is even, and } v \leq 2d - 5\}$ .

**Theorem 2.1** *Let  $(d, v) \in R$  and  $v \leq 4/5d - 17$ . Then for some  $x$ ,  $(d - x, v + x)$  and  $(x, v) \in R$ .*

**Proof** The conditions on  $d, v, x$  for  $(d - x, v + x)$  and  $(x, v) \in R$  simplify to  $x \equiv 0$  or  $3v + 1 \pmod{5}$  and  $v/2 + 1 \leq x \leq 1/3(2(d - 1) - v)$ . Since  $x$  must also be even, we are

guaranteed to find a value of  $x \equiv 0$  or  $3v + 1 \pmod{5}$  in the interval if it is at least of length ten. This is the case when  $v \leq 4/5d - 17$ .  $\square$

The alert reader may notice that  $v$  is restricted to be less than or equal to  $2d - 5$  rather than the before mentioned  $2(d - 1)$ . The reason for this will become apparent in Chapter 5.

Although Theorem 2.1 guarantees an  $x$  value for which we may use the recursion to find a design on  $(d, v)$  there are many cases in which  $v \not\leq 4/5d - 17$  and a viable  $x$  value still exists. In these cases the  $x$  value will be stated explicitly.

## Chapter 3

$a$  is even

To satisfy necessary Condition 1.1, 5 must divide  $d$  or  $d + 2v - 1$ . We will first consider when the latter case is true. This condition is satisfied when  $v$  is the maximum value,  $2(d-1)$ . Let  $a = 1/5(2(d-1) - v) \geq 0$  and this condition will also be satisfied.

**Lemma 3.1**  $(d, v) \in W$  when  $d = 2t$ ,  $v = 2(d-1) - 5a$ , and  $a$  is even.

**Proof** Let  $d = 2t$ . Let  $A = t(2t - 1 - 5/2a)$ ,  $B = 0$ ,  $\Gamma = 0$ , and  $\Delta = ta/2$ .

We require  $(a/2)$   $\delta$  base blocks. These will be constructed using *bridges*. Bridges are produced by listing the mixed differences and taking two connected arcs, the length of which are pure differences. The bridge

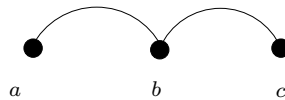


Figure 3.1: Bridge

represents the  $\delta$  base block  $((0, 1), (b, 2), (c, 2), (b - a, 1))$ . Pure differences 1 through  $a/2$  will be used exactly once in each  $\mathbb{Z}_t$  upstairs.

For example:  $(34, 16)$ ,  $a = 10$  requires 5  $\delta$  base blocks.

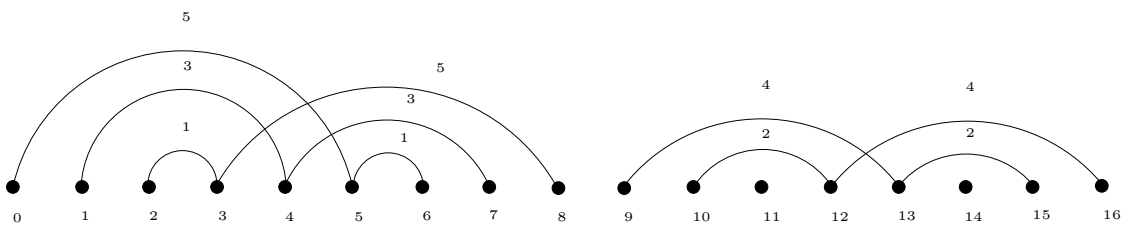


Figure 3.2: Bridges for  $(34, 16)$

The bridge

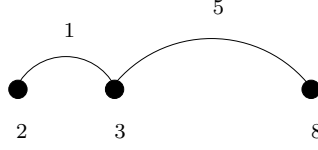


Figure 3.3: 2, 3, 8 bridge

describes the base block  $((0, 1), (3, 2), (8, 2), (1, 1))$  and uses pure differences 1 from  $\mathbb{Z}_t \times \{1\}$ , 5 from  $\mathbb{Z}_t \times \{2\}$  and mixed differences 2, 3, and 8.

The odd pure differences between 1 and  $a/2$  form the  $\delta$  base blocks:

$$((0, 1), (\lceil a/4 \rceil + j, 2), (2j + 1, 1), (3\lceil a/4 \rceil - j - 1, 2)) \text{ where } 1 \leq 2j + 1 \leq a/2$$

The even pure differences between 1 and  $a/2$  form  $\delta$  base blocks:

$$((0, 1), (3\lceil a/4 \rceil + \lfloor a/4 \rfloor + j, 2), (2j, 1), (3a/2 - j + 2, 2)) \text{ where } 2 \leq 2j \leq a/2.$$

The  $\delta$  blocks have used  $3a/2$  mixed differences and  $a/2$  pure differences. So,  $|T| = t - 3a/2$  and  $|S| = \lfloor t/2 \rfloor - a/2$ . The graph  $G[S, \{t\}, S]$  will form  $t - a$  1-factors and leaves  $t - 3a/2 - 1$  mixed difference 1-factors.

Therefore we have the required  $2t - 1 - 5/2a$   $\alpha$  base blocks. A 1-factorization of the remaining edges may be found using Lemma 1.2 to create  $\alpha$  blocks using the method described in Chapter 1.

The half difference may not be used in the construction of  $\delta$  blocks. We need  $a/2 < \lfloor d/4 \rfloor$ . However, the requirement on mixed differences is more restrictive. The  $\delta$  block construction requires the number of mixed differences, when  $a \geq 3$ , to be at least  $3a/2 + 2$ . Therefore,  $d \geq 3a + 4$ . All pairs  $d$  and  $a$  are either covered by this construction, or a special case of the recursion with values listed below.



$(38, 14) x = 10, (30, 8) x = 10, (32, 12) x = 10, (26, 10) x = 10, (24, 6) x = 4, (20, 8) x = 10,$   
 $(18, 4) x = 8, (14, 6) x = 4, (12, 2) x = 2$

□

Chapter 4

$a$  is odd

**Lemma 4.1**  $(d, v) \in W$  when  $t = 3s$  and  $v = 2(d - 1) - 5a$ .

**Proof** The cases where  $v \leq 3$  have been previously solved. Lemma 3.1 solves the case when  $a$  is even. We now assume that  $v > 3$  and  $a$  is odd.

Let  $A = t(2t - 1 - 4a + 3(a - 1)/2)$ ,  $B = 2t$ ,  $\Gamma = 0$ ,  $\Delta = t(a - 1)/2$ .

The  $\beta$  blocks are constructed using pure differences 1 and 2 and the first 3 vertices in  $V$ . Each  $\beta$  connects a vertex in  $V$  with 3 vertices in  $\mathbb{Z}_t$ . The vertices in  $\mathbb{Z}_t$  can be partitioned into three parallel classes. Each class contains  $s$  disjoint sets of size 3. These sets of size 3, along with a single vertex in  $V$ , will create a  $\beta$  block using an edge of pure difference 1 and 2. Each of the three classes uses  $s$  edges of each difference 1 and 2 therefore exhausting all  $2t$  edges of these differences.

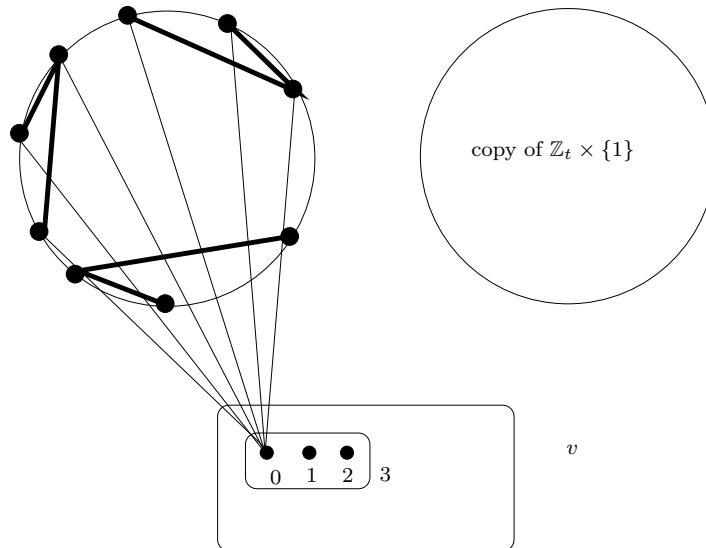


Figure 4.1: One class of  $\beta$  blocks

All  $\beta$  blocks are defined by:

$$((i, 3), (i+3j, 1), (i+3j+2, 1), (i+3j+1, 1)) \text{ and } ((i, 3), (i+3j, 2), (i+3j+2, 2), (i+3j+1, 2))$$

where  $0 \leq i \leq 2$  and  $0 \leq j \leq s-1$ .

The  $(a-1)/2$   $\delta$  base blocks are constructed using a similar method as in Lemma 3.1. In this case however, since pure differences 1 and 2 have been used in the  $\beta$  blocks, we use pure differences 3 through  $(a-1)/2+2$ .

$$\text{Let } \delta' = (a-1)/2.$$

The following are  $\delta$  base blocks and are developed (mod  $t$ ):

$$((0, 1), (\lceil \delta'/2 \rceil + 1 + j, 2), (2j+1, 1), (3\lceil \delta'/2 \rceil + 4 - j, 2)) \text{ where } 3 \leq 2j+1 \leq \delta' + 2$$

$$((0, 1), (3\lceil \delta'/2 \rceil + \lfloor \delta'/2 \rfloor + 5 + j, 2), (2j, 1), (3\delta' + 11 - j, 2)) \text{ where } 4 \leq 2j \leq \delta' + 2$$

The half difference may not be used in the construction of  $\delta$  blocks. So, we need

$(a-1)/2+2 < \lfloor d/4 \rfloor$ . However, the requirement on mixed differences is more restrictive.

The  $\delta$  block construction requires the number of mixed differences, when  $a \geq 4$ , ( $\delta' \geq 2$ ), to be at least  $3\delta' + 10$ . Therefore,  $d \geq 3a + 17$ .

The shaded region in the following graph, referred to as the dead zone, indicates pairs  $(a, d)$  that are not covered by this construction and are not guaranteed to be covered by the recursion.

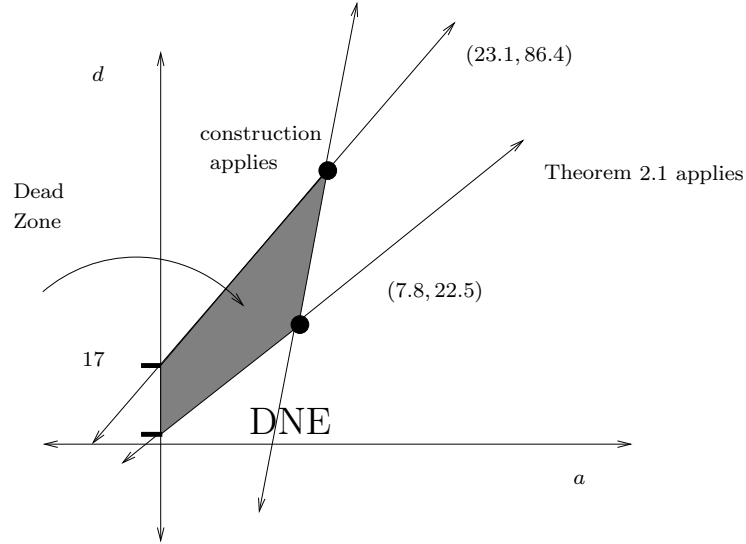


Figure 4.2: Dead Zone for Lemma 4.1

Several of these cases can be covered by modifying the above the construction to using only odd pure differences in the  $\delta$  blocks by replacing the  $\delta$  base blocks with:

$$((0, 1), (\delta' + 1 + j, 2), (2j + 1, 1), (3\delta' + 4 - j, 2)) \text{ where } 1 \leq j \leq \delta'$$

This modification is used for designs on:  $(54, 41)$ ,  $(48, 39)$ ,  $(42, 37)$ ,  $(36, 35)$ ,  $(30, 33)$ ,  $(30, 23)$  and  $(24, 21)$ . The remaining dead zone cases needed for this lemma are:

$(60, 43)$ : In the main construction above replace  $\delta$  base blocks with:

$$((0, 1), (8 + j, 2), (2j + 1, 1), (25 - j, 2)) \text{ where } 1 \leq j \leq 6 \text{ and } ((0, 1), (14, 2), (8, 1), (22, 2))$$

$(42, 27)$ : In the main construction above replace the  $\delta$  base blocks with:

$$((0, 1), (5 + j, 2), (2j + 1, 1), (16 - j, 2)) \text{ where } 3 \leq 2j + 1 \leq 9$$

$$((0, 1), (11, 2), (6, 1), (17, 2))$$

$(36, 25)$ : In the main construction above replace the  $\delta$  base blocks with:

$$((0, 1), (4 + j, 2), (2j + 1, 1), (13 - j, 2)) \text{ where } 3 \leq 2j + 1 \leq 7$$

$$((0, 1), (9, 2), (6, 1), (15, 2))$$

$(12, 7)$ :  $\beta$  blocks:  $((0, 3), (5, 1), (2, 1), (1, 1)), ((0, 3), (3, 2), (3, 1), (1, 2)), ((0, 3), (0, 1),$

$(4, 1), (0, 2), ((0, 3), (2, 2), (5, 2), (4, 2)), ((1, 3), (3, 1), (0, 1), (5, 1)),$   
 $((1, 3), (1, 2), (1, 1), (5, 2)), ((1, 3), (4, 1), (2, 1), (4, 2)), ((1, 3), (0, 2), (3, 2), (2, 2)),$   
 $((2, 3), (1, 1), (4, 1), (3, 1)), ((2, 3), (2, 1), (0, 1), (2, 2)), ((2, 3), (5, 2), (5, 1), (3, 2)),$   
 $((2, 3), (4, 2), (0, 2), (1, 2))$

$\delta$  base block:  $((0, 1), (2, 2), (1, 1), (3, 2))$

Create  $\alpha$  blocks with pure differences 4 and 5 and the remaining four vertices in  $V$ .

$(6, 5)$ : Let  $\{0, 1, 2, 3, 4, 5\}$  be the vertices of  $K_6$ .

$(1, (0, 3), 5, 0), (4, (0, 3), 2, 3), (2, (1, 3), 0, 1), (5, (1, 3), 3, 4), (3, (2, 3), 1, 2),$   
 $(0, (2, 3), 4, 5), (0, 3, (3, 3), (4, 3)), (1, 4, (3, 3), (4, 3)), (2, 5, (3, 3), (4, 3)).$

Special cases of the Recursion:  $(84, 51) x = 30, (78, 49) x = 28, (72, 47) x = 30, (66, 45) x = 26,$   
 $(60, 33) x = 20, (54, 31) x = 20, (48, 29) x = 18, (42, 17) x = 20, (36, 15) x = 10,$   
 $(30, 13) x = 10, (24, 11) x = 10, (18, 9), x = 8. \quad \square$

**Lemma 4.2**  $(d, v) \in W$  when  $t = 3s + 1$  and  $v = 2(d - 1) - 5a$ .

**Proof** The cases where  $v \leq 3$  have been previously solved. Lemma 3.1 solves the case when  $a$  is even. We now assume that  $v > 3$  and  $a$  is odd.

Let  $A = t(2t - 1 - 4a + 3(a - 1))/2, B = 2t, \Gamma = 0, \Delta = t(a - 1)/2$ . Similar to Lemma 4.1,  $s$   $\beta$  blocks are constructed using pure differences 1 and 2 with each of the first 3 vertices in  $V$ . We need one more  $\beta$ , which will contain vertex  $(3, 3)$ , to exhaust the edges of pure differences 1 and 2. Mixed difference 0 will then be used to create  $\alpha$  blocks with the first five vertices in  $V$ . These  $\alpha$  blocks will use all remaining edges adjacent to those five vertices in  $V$ . Here are the details:

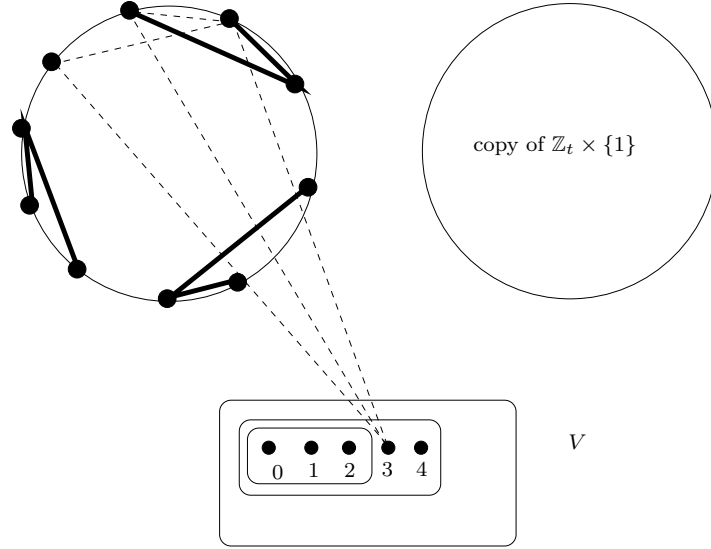


Figure 4.3: Extra  $\beta$  block needed in Lemma 4.2

All  $\beta$  blocks are defined by:

$((i, 3), (i+3j+2, 1), (i+3j, 1), (i+3j+1, 1)), ((i, 3), (i+3j+2, 2), (i+3j, 2), (i+3j+1, 2))$  where  $0 \leq i \leq 2$  and  $0 \leq j \leq s-1$ .  $((3, 3), (1, 1), (0, 1), (t-1, 1))$  and  $((3, 3), (1, 2), (0, 2), (t-1, 2))$ . Since mixed difference 0 is used in the  $\alpha$  blocks, the  $\delta$  bridges begin at mixed difference 1.

Let  $\delta' = (a-1)/2$ .

The following are  $\delta$  base blocks and are developed (mod  $t$ ):

$((0, 1), (\lceil \delta'/2 \rceil + 2 + j, 2), (2j+1, 1), (3\lceil \delta'/2 \rceil + 5 - j, 2))$  where  $3 \leq 2j+1 \leq \delta' + 2$

$((0, 1), (3\lceil \delta'/2 \rceil + \lfloor \delta'/2 \rfloor + 6 + j, 2), (2j, 1), (3\delta' + 12 - j, 2))$  where  $4 \leq 2j \leq \delta' + 2$

The last blocks to be constructed are  $\alpha$  blocks. First, the remaining edges between the vertices upstairs and the first five vertices in  $V$  will be exhausted. Each vertex upstairs needs to be connected with two of the five vertices downstairs.

These  $\alpha$  blocks are as follows:

$((1, 1), (1, 2), (2, 3), (4, 3)), ((0, 1), (0, 2), (1, 3), (4, 3)), ((t-1, 1), (t-1, 2), (0, 3), (4, 3)), ((t-2, 1), (t-2, 2), (3, 3), (4, 3))$  and  $((j, 1), (j, 2), (3, 3), (4, 3))$  where  $2 \leq j \leq t-3$

All other  $\alpha$  blocks are constructed as described in Chapter one with all vertices  $v \geq 5$ .

The dead zone indicates pairs  $(a, d)$  that are not covered by this construction and are not guaranteed by Theorem 2.1.

We must also ensure that all cases when  $v \leq 5$  are covered.

As in Lemma 4.1, many of these cases are covered by the following modification on  $\delta$  blocks where only odd pure differences are used.

Use the above construction replacing the  $\delta$  base blocks with:

$$((0, 1), (\delta' + 2 + j, 2), (2j + 1, 1), (3\delta' + 5 - j, 2)) \text{ where } 1 \leq j \leq \delta'$$

This modification works for the following cases:  $(70, 43), (64, 41), (62, 47), (56, 45), (50, 43), (44, 31), (38, 29), (38, 39), (32, 27), (32, 37),$  and  $(26, 25)$ .

$(14, 11)$ :

$\delta$  base block:  $((0, 1), (0, 2), (3, 1), (3, 2))$

$\beta$  blocks:  $((0, 1), (j, 3), (5, 1), (6, 1)), ((0, 2), (j, 3), (5, 2), (6, 2))$  where  $1 \leq j \leq 6$

$\alpha$  blocks:  $((1, 1), (2, 2), (0, 3), (5, 3)), ((2, 1), (1, 2), (0, 3), (5, 3)), ((3, 1), (4, 2), (0, 3), (2, 3)),$

$((2, 1), (3, 2), (1, 3), (6, 3)), ((3, 1), (1, 2), (1, 3), (6, 3)), ((4, 1), (5, 2), (1, 3), (3, 3)),$

$((5, 1), (4, 2), (1, 3), (3, 3)), ((5, 1), (6, 2), (4, 3), (2, 3)), ((6, 1), (5, 2), (4, 3), (2, 3)),$

$((6, 1), (0, 2), (5, 3), (3, 3)), ((0, 1), (6, 2), (5, 3), (3, 3)), ((0, 1), (1, 2), (6, 3), (4, 3)),$

$((1, 1), (0, 2), (6, 3), (4, 3))$ . The two one factors from mixed differences 2 and 5 are then used to make  $\alpha$  blocks with vertices 7, 8, 9, and 10 in  $V$ .

$(20, 13)$ : In the main construction above replace  $\delta$  base blocks with:  $((0, 1), (9, 2), (4, 1), (3, 2))$  and  $((0, 1), (4, 2), (3, 1), (7, 2))$ .

Special cases of the Recursion:

$(92, 57) \ x = 40, (86, 55) \ x = 30, (80, 53) \ x = 30, (74, 51) \ x = 30, (68, 49) \ x = 28, (68, 39) \ x = 30, (62, 37) \ x = 22, (56, 35) \ x = 20, (50, 33) \ x = 20, (44, 21) \ x = 20, (38, 19) \ x = 18, (32, 17) \ x = 12, (26, 15) \ x = 10, (26, 5) \ x = 10.$  □

**Lemma 4.3**  $(d, v) \in W$  when  $t = 3s + 2$  and  $v = 2(d - 1) - 5a$ .

**Proof** The cases where  $v \leq 3$  have been previously solved. Lemma 3.1 solves the case when  $a$  is even. We now assume that  $v > 3$  and  $a$  is odd.

$$\text{Let } A = t(2t - 1 - 4a + 3(a - 1)/2), B = 2t, \Gamma = 0, \Delta = t(a - 1)/2$$

Similar to Lemma 4.1, in each  $\mathbb{Z}_t \times \{1\}$  and  $\mathbb{Z}_t \times \{2\}$ ,  $s$   $\beta$  blocks are constructed using pure differences 1 and 2 with the first 3 vertices in  $v$ . We need two more  $\beta$ 's in each  $\mathbb{Z}_t \times \{1\}$  and  $\mathbb{Z}_t \times \{2\}$  to exhaust the edges of pure differences 1 and 2. One of which will contain vertiex  $(3, 3)$ , and the other  $(4, 3)$ . Mixed difference 0 will then be used to create  $\alpha$  blocks with the first five vertices in  $V$ . These  $\alpha$  blocks will use all remaining edges adjacent to those first five vertices in  $V$ . Here are the details:

All  $\beta$  blocks are defined by:

$$\begin{aligned} &((i, 3), (i+3j+2, 1), (i+3j, 1), (i+3j+1, 1)), ((i, 3), (i+3j+2, 2), (i+3j, 2), (i+3j+1, 2)) \text{ where} \\ &0 \leq i \leq 2 \text{ and } 0 \leq j \leq s - 1. ((4, 3), (1, 1), (0, 1), (t - 1, 1)), ((4, 3), (1, 2), (0, 2), (t - 1, 2)), \\ &((3, 3), (0, 1), (t - 1, 1), (t - 2, 1)), \text{ and } ((3, 3), (0, 2), (t - 1, 2), (t - 2, 2)) \end{aligned}$$

Since mixed difference 0 is used in the  $\alpha$  blocks, the  $\delta$  bridges begin at mixed difference 1.

$$\text{Let } \delta' = (a - 1)/2.$$

The following are  $\delta$  base blocks and are developed (mod  $t$ ):

$$\begin{aligned} &((0, 1), (\lceil \delta'/2 \rceil + 2 + j, 2), (2j + 1, 1), (3\lceil \delta'/2 \rceil + 5 - j, 2)) \text{ where } 3 \leq 2j + 1 \leq \delta' + 2 \\ &((0, 1), (3\lceil \delta'/2 \rceil + \lfloor \delta'/2 \rfloor + 6 + j, 2), (2j, 1), (3\delta' + 12 - j, 2)) \text{ where } 4 \leq 2j \leq \delta' + 2 \end{aligned}$$

Now the remaining edges between the vertices upstairs and the first five vertices in  $V$  will be exhausted. Each vertex upstairs will be connected with two of the first five vertices downstairs in  $\alpha$  blocks as follows:

$$\begin{aligned} &((0, 1), (0, 2), (2, 3), (1, 3)), ((1, 1), (1, 2), (2, 3), (3, 3)), ((t - 1, 1), (t - 1, 2), (1, 3), \\ &(0, 3)), ((t - 2, 1), (t - 2, 2), (3, 3), (1, 3)), ((t - 3, 1), (t - 3, 2), (3, 3), (4, 3)) \text{ and} \\ &((j, 1), (j, 2), (3, 3), (4, 3)) \text{ where } 2 \leq j \leq t - 4. \text{ All other } \alpha \text{ blocks are constructed as described} \\ &\text{in Chapter one with the remaining vertices } v \geq 5. \end{aligned}$$



This lemma has the same restrictions as Lemma 4.2 and thus has the same dead zone. The same modification on the  $\delta$  blocks as in Lemmas 4.1 and 4.2 cover the following dead zone cases:  $(46, 35)$ ,  $(34, 31)$ ,  $(40, 33)$ ,  $(62, 47)$ ,  $(28, 29)$ , and  $(22, 17)$ .

$(52, 37)$ : Use the main construction for  $\alpha$  and  $\beta$  blocks. Use the following  $\delta$  base blocks:  $((0, 1), (7 + j, 2), (2j + 1, 1), (20 - j, 2))$  where  $1 \leq j \leq 5$  and  $((0, 1), (14, 2), (8, 1), (22, 2))$ .

$(58, 39)$ : Use the main construction of this lemma for  $\beta$  and  $\alpha$  blocks. The  $\delta$  base blocks are  $((0, 1), (15 - j, 2), (15 - 2j, 1), (16 + j, 2))$  where  $3 \leq 2j + 1 \leq 13$  and  $((0, 1), (16, 2), (8, 1), (24, 2))$ .

$(40, 23)$ : Use main construction above replacing the  $\delta$  base blocks with:  $((0, 1), (6, +j, 2), (2j + 1, 1), (17 - j, 2))$  where  $1 \leq j \leq 4$  and  $((0, 1), (12, 2), (6, 1), (18, 2))$ .

$(28, 19)$ : Use the main construction of this lemma for  $\beta$  and  $\alpha$  blocks. The  $\delta$  base blocks are  $((0, 1), (5, 2), (3, 1), (10, 2))$ ,  $((0, 1), (6, 2), (5, 1), (9, 2))$ , and  $((0, 1), (4, 2), (8, 1), (12, 2))$ .

Special cases of the Recursion:  $(76, 45)$   $x = 30$ ,  $(52, 27)$   $x = 20$ ,  $(46, 25)$   $x = 16$ ,  $(34, 21)$   $x = 14$ ,  $(34, 11)$   $x = 10$ ,  $(28, 9)$   $x = 8$ ,  $(22, 7)$   $x = 10$ ,  $(16, 5)$   $x = 6$ . □

Chapter 5

$$5|d$$

**Lemma 5.1** *When  $d$  is even,  $5|d$ , and  $v = 2d - 3$  or  $v = 2d - 4$  a  $K_4 - e$  design does not exist, on  $K_d + v$ .*

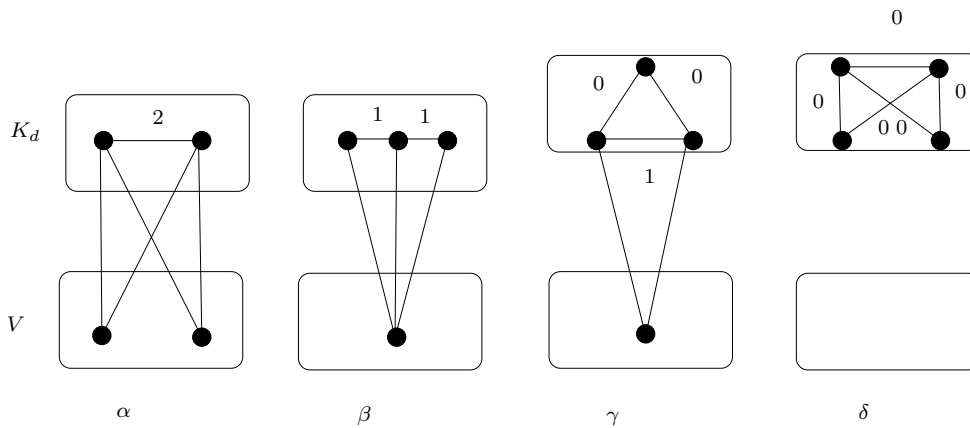


Figure 5.1: Number of edge colors

**Proof** Assign each vertex in  $V$  a unique color. We then assign colors to the edges upstairs in accordance with the block type and vertices downstairs as follows: Each edge upstairs in an  $\alpha$  block will be assigned two colors. These are the colors of the two vertices in  $V$  that are contained in the  $\alpha$  block. Each edge upstairs in a  $\beta$  block will be assigned the one color of the vertex  $v$  contained in the  $\beta$  block. Let  $e$  be the edge in a  $\gamma$  block that is incident to both vertices upstairs that are adjacent to the vertex in  $v$ . Color  $e$  the same color as the vertex in  $V$ . The other two edges in a  $\gamma$  block and all edges in a  $\delta$  block will not be assigned colors.

Let  $p$  be the number of edges upstairs that are assigned two colors. Let  $q$  be the number of edges upstairs that are assigned one color in  $\gamma$  blocks. Let  $r$  be the number of pairs of one colored edges in  $\beta$  blocks and  $s$  is the number of uncolored edges upstairs.

Counting edges upstairs at a particular vertex we get the equation  $d - 1 = p + q + 2r + s$ . Counting colors at each a vertex upstairs we get  $v = 2p + q + r$ . From necessary Condition 1.2 we know that  $2(d - 1) - v \geq 0$ . Call this value the *deficiency*.

The deficiency,  $2(d - 1) - v = q + 3r + 2s$ .

**Case 1:** The deficiency =  $1 = q + 3r + 2s$ .

Thus,  $q = 1, r = s = 0$ . Since this is the only possible combination, all vertices upstairs must have this coloring. However, the single one colored edge ( $q = 1$ ) cannot come from a  $\gamma$  block because there are no uncolored edges ( $s = 0$ ). The single one colored edge cannot come from a  $\beta$  block because there are no pairs of single colored edges ( $r = 0$ ). The other block types do not contain single one colored edges. Therefore, the deficiency cannot equal one and there does not exist a  $K_4 - e$  design when  $v = 2d - 3$ .

**Case 2:** The deficiency =  $2 = q + 3r + 2s$ .

There are two possibilities,  $2 = q$  and  $r = s = 0$  or  $q = r = 0$  and  $s = 1$ . In either case, since  $r = 0$ , there are no  $\beta$  blocks. The two single one colored edges ( $q = 2$ ) must come from  $\gamma$  blocks. However, a  $\gamma$  block cannot exist unless there is a vertex with 2 uncolored edges and  $s = 0$  or 1. Therefore, the deficiency cannot equal two and there does not exist a  $K_4 - e$  design when  $v = 2d - 4$ . □

**Lemma 5.2** *If  $(d, v) \in W$  then  $(dk, v + 2d(k - 1)) \in W$ .*

**Proof** Create the graph  $K_{dk} + v$  by taking each vertex upstairs in  $K_d + v$  with a  $(d, v)$  design and blowing it up by  $k$ . Then  $k$  copies of the  $(d, v)$  design will exhaust all edges incident to vertices in  $V$  and  $(\frac{kd(d-1)}{2})$  edges upstairs. By Lemma 1.2 the remaining edges upstairs can be partitioned into  $d(k - 1)$  one-factors. For each one-factor, add a pair of vertices downstairs and the appropriate edges to obtain  $K_{dk} + (v + 2d(k - 1))$ . Now each one-factor and pair of vertices downstairs can be partitioned into  $\alpha$  blocks. □

$K_{a,b,c}$  is a complete multipartite graph. Let  $S = \{(a, b, c): \text{there exists a } K_4 - e \text{ design on } K_{a,b,c}\}$ .

**Lemma 5.3** *If  $(a, b, c) \in S$ ,  $(b, v)$  and  $(c, v) \in W$  then  $(b + c, a + v) \in W$ .*

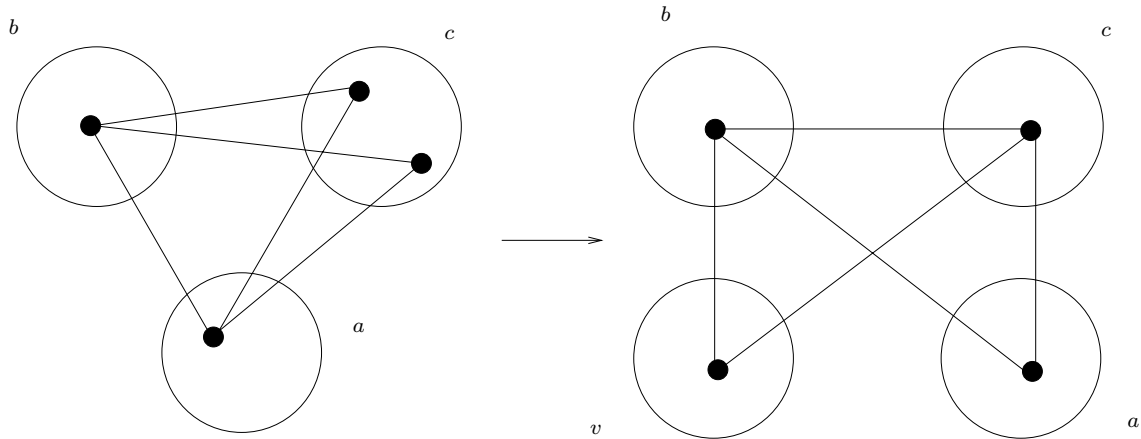


Figure 5.2: Another construction

**Proof** The union of the blocks from the designs  $(a, b, c)$ ,  $(b, v)$ , and  $(c, v)$  create a  $(b + c, a + v)$  design. □

**Lemma 5.4**  $(10, 10, 10) \in S$ .

**Proof** The following blocks form  $(10, 10, 10)$ :

$$((0 + i, a), (0 + i, b), (2 + i, c), (3 + i, c))$$

$$((0 + i, b), (0 + i, c), (2 + i, a), (3 + i, a))$$

$$((0 + i, c), (0 + i, a), (2 + i, b), (3 + i, b))$$

$$((5 + i, a), (0 + i, b), (1 + i, c), (6 + i, c))$$

$$((5 + i, b), (0 + i, c), (1 + i, a), (6 + i, a))$$

$$((5 + i, c), (0 + i, a), (1 + i, b), (6 + i, b)), \text{ for } 0 \leq i \leq 9$$

□

**Lemma 5.5**  $(d, v) \in W$  when  $5|d$ , and  $v \leq 2(d-1)$  and  $v \neq 2d-3$ ,  $v \neq 2d-4$ .

**Proof** Since  $5|d$  and  $d$  is even, let  $d = 10t$ .

**Case 1:**  $t = 1$

Previously solved cases:  $v = 1, 2, 3$  [6],  $v = 4$  [6, example 4.2],  $v = 7$  [6, example 4.5], and  $v = 9$  [6, example 4.9]

When  $v = 8, 13$  and  $18$ ,  $v = 2(d-1) - 5a$  and are solved in Chapters 3 and 4.

(10, 4): Let  $\{1, 2, 3, 4, 5, 6, 7, 8, 9, 10\}$  be the vertices upstairs  $(5, 6, (2, 3), (3, 3)), (4, 10, (2, 3), (3, 3)), (3, 9, (2, 3), (3, 3)), (2, 8, (2, 3), (3, 3)), (1, 7, (2, 3), (3, 3)), ((1, 3), 3, 1, 7), (1, 4, 2, 8), (2, 5, 3, 9), (3, (0, 3), 4, 10), (4, 7, 5, 6), (5, 8, (0, 3), (1, 3)), ((0, 3), 9, 7, 1), (7, 10, 8, 2), (8, 6, 9, 3), (9, (1, 3), 10, 4), (10, 1, 6, 5), (6, 2, (1, 3), (0, 3))$ .

(10, 5): Let  $\{0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}$  be the set of vertices upstairs.  $((1, 3), 1, 0, 2), (9, 3, (1, 3), 1), ((1, 3), 6, 5, 7), (8, 3, (1, 3), 6), (7, 1, 5, 8), (7, 2, 3, 4), (6, 2, 0, 3), (3, 4, 5, 0), (8, 9, 5, 0), (1, 4, (2, 3), (3, 3)), (6, 9, (2, 3), (3, 3)), (6, 1, (0, 3), (4, 3)), (0, 4, (0, 3), (4, 3)), (5, 2), (2, 3), (3, 3)), (8, 3, (2, 3), (3, 3)), (7, 0, (2, 3), (3, 3)), (5, 0, (0, 3), (4, 3)), (7, 3, (0, 3), (4, 3)), (8, 2, (0, 3), (4, 3)). [7, page 12]$

(10, 6):  $\mathbb{Z}_5 \times \{1\} \cup (0, 3)$  and  $\mathbb{Z}_5 \times \{2\} \cup (0, 3)$  each form a copy of  $K_6$  and can be decomposed in 6 blocks. The following blocks complete the design.  $((1, 3), (0, 1), (0, 2), (1, 2)), ((2, 3), (1, 1), (1, 2), (2, 2)), ((3, 3), (2, 1), (2, 2), (3, 2)), ((4, 3), (3, 1), (3, 2), (4, 2)), ((5, 3), (4, 1), (4, 2), (0, 2)), ((1, 3), (2, 2), (4, 1), (3, 1)), ((2, 3), (3, 2), (0, 1), (4, 1)), ((3, 3), (4, 2), (1, 1), (0, 1)), ((4, 3), (0, 2), (2, 1), (1, 1)), ((5, 3), (1, 2), (3, 1), (2, 1)), ((0, 1), (2, 2), (4, 3), (5, 3)), ((1, 1), (3, 2), (1, 3), (5, 3)), ((2, 1), (4, 2), (1, 3), (2, 3)), ((3, 1), (0, 2), (2, 3), (3, 3)), ((4, 1), (1, 2), (3, 3), (4, 3))$ .

(10, 7): Let  $\{1, 2, 3, 4, 5, 6, 7, 8, 9, 10\}$  be the vertices upstairs.  $((0, 3), 7, 5, 8), ((0, 3), 10, 6, 9), (1, 2, (0, 3), (4, 3)), (3, 4, (0, 3), (1, 3)), ((1, 3), 1, 7, 9), ((1, 3), 2, 10, 8), (6, 5, (1, 3), (2, 3)), ((2, 3), 8, 1, 3), ((2, 3), 9, 2, 4), (7, 10, (2, 3), (3, 3)), ((3, 3), 4, 1, 5), ((3, 3), 3, 2, 6), (8, 9, (3, 3), (4, 3)), ((4, 3), 5, 3, 10), ((4, 3), 6, 4, 7), ((5, 3), 5, 1, 8), (3, 10, (5, 3), 1), ((5, 3), 6, 2, 9), (4, 7, (5, 3), 2), ((6, 3), 5, 2, 9), (3, 7, (6, 3), 9), ((6, 3), 6, 1, 8), (4, 10, (6, 3), 8)$ .

(10, 9): Let  $\{1, 2, 3, 4, 5, 6, 7, 8, 9, 10\}$  be the vertices upstairs.  $((0, 3), 1, 4, 10)$ ,  
 $((0, 3), 7, 2, 3)$ ,  $(5, 8, (0, 3), (2, 3))$ ,  $(6, 9, (0, 3), (1, 3))$ ,  $((1, 3), 1, 2, 8)$ ,  
 $((1, 3), 5, 3, 4)$ ,  $(7, 10, (1, 3), (2, 3))$ ,  $((2, 3), 1, 3, 9)$ ,  $((2, 3), 6, 2, 4)$ ,  $(1, 5, (3, 3), (4, 3))$ ,  
 $(2, 8, (3, 3), (4, 3))$ ,  $(3, 4, (3, 3), (4, 3))$ ,  $(6, 7, (3, 3), (4, 3))$ ,  $(9, 10, (3, 3), (4, 3))$ ,  
 $(1, 6, (5, 3), (6, 3))$ ,  $(2, 4, (5, 3), (6, 3))$ ,  $(3, 9, (5, 3), (6, 3))$ ,  $(5, 7, (5, 3), (6, 3))$ ,  
 $(8, 10, (5, 3), (6, 3))$ ,  $(1, 7, (7, 3), (8, 3))$ ,  $(2, 3, (7, 3), (8, 3))$ ,  $(4, 10, (7, 3), (8, 3))$ ,  
 $(8, 9, (7, 3), (8, 3))$ ,  $(2, 5, 9, 10)$ ,  $(3, 6, 8, 10)$ ,  $(4, 7, 8, 9)$ .

(10, 10): Let  $\{0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}$  be the set of vertices upstairs.  
 $((i, 3), i, 8, 9)$  for  $0 \leq i \leq 7$ ,  $((0, 3), 7, 2, 1)$ ,  $((6, 3), 5, 0, 7)$ ,  $((7, 3), 6, 1, 0)$ ,  $(0, 1, (4, 3)(6, 3))$ ,  
 $(1, 2, (5, 3), (7, 3))$ ,  $(2, 2, (0, 3), (6, 3))$ ,  $(3, 4, (1, 3), (7, 3))$ ,  $(4, 5, (0, 3), (2, 3))$ ,  
 $(5, 6, (1, 3), (3, 3))$ ,  $(6, 7, (2, 3), (4, 3))$ ,  $(7, 0, (5, 3)(3, 3))$ ,  
 $(8, 9, (8, 3), (9, 3))$ ,  $(0, 4, (8, 3), (9, 3))$ ,  $(1, 5, (8, 3), (9, 3))$ ,  
 $(2, 6, (8, 3), (9, 3))$ ,  $(3, 7, (8, 3), (9, 3))$ .

(10, 11): Put a  $(5, 1)$  design on  $\mathbb{Z}_5 \times \{1\} \cup (0, 3)$  and  $\mathbb{Z}_5 \times \{2\} \cup (0, 3)$ . Then use mixed differences 0 through 4 to create  $\alpha$  blocks with the last 10 vertices in  $V$ .

(10, 12): Let  $\{1, 2, 3, 4, 5, 6, 7, 8, 9, 10\}$  be the vertices upstairs.  $((i, 3), i, 7, 8)$  for  $1 \leq i \leq 6$ ,  
 $((j + 1, 3), j, 9, 10)$  for  $1 \leq j \leq 5$ ,  $((1, 3), 6, 9, 10)$ ,  
 $(1, 6, (3, 3), (4, 3))$ ,  $(1, 2, (5, 3), (6, 3))$ ,  $(3, 6, (2, 3), (3, 3))$ ,  $(3, 4, (1, 3), (6, 3))$ ,  
 $(1, 5, (7, 3), (8, 3))$ ,  $(4, 6, (7, 3), (8, 3))$ ,  $(2, 3, (7, 3), (8, 3))$ ,  $(7, 8, (7, 3), (8, 3))$ ,  
 $(9, 10, (7, 3), (8, 3))$ ,  $(1, 3, (11, 3), (12, 3))$ ,  $(2, 4, (11, 3), (12, 3))$ ,  $(5, 6, (11, 3), (12, 3))$ ,  
 $(7, 9, (11, 3), (12, 3))$ ,  $(8, 10, (11, 3), (12, 3))$ ,  $(1, 4, (9, 3), (10, 3))$ ,  $(2, 6, (9, 3), (10, 3))$ ,  
 $(3, 5, (9, 3), (10, 3))$ ,  $(8, 9, (9, 3), (10, 3))$ ,  $(7, 10, (9, 3), (10, 3))$ .

(10, 14):  $((1, 3), (0, 1), (4, 1), (1, 1))$ ,  $((1, 3), (0, 2), (4, 2), (1, 2))$ ,  
 $((2, 3), (1, 1), (4, 1), (2, 1))$ ,  $((2, 3), (1, 2), (4, 2), (2, 2))$ ,  $((3, 3), (2, 1), (4, 1), (3, 1))$ ,  
 $((3, 3), (2, 2), (4, 2), (3, 2))$ ,  $((4, 3), (3, 1), (4, 1), (0, 2))$ ,  $((4, 3), (3, 3), (4, 2), (0, 1))$ ,  
 $((3, 1), (3, 2), (1, 3), (2, 3))$ ,  $((2, 1), (2, 2), (1, 3), (4, 3))$ ,  $((1, 1), (1, 2), (3, 3), (4, 3))$ ,  
 $((0, 1), (0, 2), (2, 3), (3, 3))$ ,  $((3, 1), (4, 2), (5, 3), (6, 3))$ ,  $((2, 1), (0, 1), (4, 3), (6, 3))$ ,

$((0, 1), (0, 2), (5, 3), (6, 3)), ((4, 1), (2, 2), (5, 3), (6, 3)), ((3, 2), (1, 2), (5, 3), (6, 3)),$   
 $((3, 1), (2, 2), (7, 3), (8, 3)), ((2, 1), (0, 2), (7, 3), (8, 3)), ((1, 1), (3, 2), (7, 3), (8, 3)),$   
 $((0, 1), (1, 2), (7, 3), (8, 3)), ((4, 2), (4, 1), (7, 3), (8, 3)), ((3, 1), (0, 1), (9, 3), (10, 3)),$   
 $((2, 1), (4, 2), (9, 3), (10, 3)), ((1, 1), (2, 2), (9, 3), (10, 3)), ((4, 1), (1, 2), (9, 3), (10, 3)),$   
 $((3, 2), (1, 2), (0, 3), (10, 3)), ((3, 1), (1, 2), (11, 3), (12, 3)), ((2, 1), (3, 2), (11, 3), (12, 3)),$   
 $((1, 4), (4, 2), (11, 3), (12, 3)), ((0, 1), (2, 2), (11, 3), (12, 3)), ((4, 1), (0, 2), (11, 3), (12, 3)),$   
 $((3, 1), (1, 1), (13, 3), (14, 3)), ((0, 1), (4, 2), (13, 3), (14, 3)), ((4, 1), (3, 2), (13, 3), (14, 3)),$   
 $((0, 1), (4, 2), (13, 3), (14, 3)), ((4, 1), (3, 2), (13, 3), (14, 3)), ((2, 2), (0, 2), (13, 3), (14, 3)),$   
 $((2, 1), (1, 2), (13, 3), (14, 3)).$

$(10, 15): ((1, 3), (4, 1), (0, 1), (1, 1)), ((1, 3), (4, 1), (0, 2), (1, 1)),$   
 $((2, 3), (3, 1), (0, 1), (1, 1)), ((2, 3), (3, 2), (0, 2), (1, 2)), ((3, 3), (2, 1), (0, 1), (1, 1)),$   
 $((3, 3), (2, 2), (0, 2), (1, 2)), ((0, 1), (0, 2), (6, 3), (7, 3)), ((1, 1), (1, 2), (6, 3), (7, 3)),$   
 $((2, 1), (2, 2), (4, 3), (7, 3)), ((3, 1), (3, 2), (6, 3), (7, 3)), ((4, 1), (4, 2), (5, 3), (7, 3)),$   
 $((3, 1)(2, 1), (1, 3), (5, 3)), ((3, 2), (2, 2), (1, 3), (5, 3)), ((4, 1), (2, 1), (2, 3), (6, 3)),$   
 $((4, 2), (2, 2), (2, 3), (6, 3)), ((3, 1), (4, 1), (3, 3), (4, 3)), ((3, 2), (4, 2), (3, 3), (4, 3)),$   
 $((0, 1), (1, 1), (4, 3), (5, 3)), ((0, 2), (1, 2), (4, 3), (5, 3)).$  Use mixed differences 1, 2, 3, and 4 to make  $\alpha$  blocks with vertices 8 through 15 in  $V$ .

We will now assume that  $t > 1$ .

(The edges of pure difference  $t$  and  $2t$  and mixed differences  $0, t, 2t, 3t$ , and  $4t$  produce  $t$  copies of  $K_{10}$  upstairs. This set of edges is replaced with  $K_{10} + h$  in Lemma 5.2).

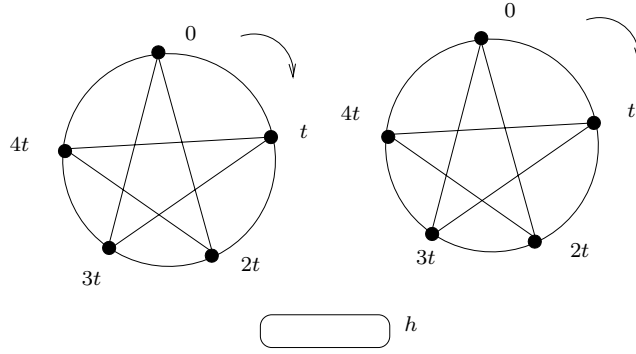


Figure 5.3:  $K_{10t} + h$

For  $20t - 20 \leq v \leq 20t - 2$ , use  $(10, h)$ ,  $0 \leq h \leq 18$ , in Lemma 5.2. In these designs,  $d(k - 1)$  one-factors are used to make  $\alpha$  blocks with  $2d(k - 1)$  vertices downstairs. For every 5 sets of one-factors that we can turn into  $\delta$  blocks, we reduce the number of vertices downstairs by 10. For each  $\delta$  base block we can construct, we can cover the values of  $v$  where  $20t - 20 - 10\delta \leq v \leq 20t - 2 - 10\delta$ . Here are the details:

**Case 2:**  $t$  is odd

$\delta$  base blocks:

$((0, 1), (2t - 1 - i, 2), (2t - 2 - 2i, 1), (2t + 1 + i, 2))$  where  $0 \leq i \leq t - 2$

$((0, 1), (3t + 2, 2), (1, 1), (3t + 3, 2))$ ,

$((0, 1), (4t + 1 + j, 2), (3 + 2j, 1), (5t - 1 - j, 2))$  where  $0 \leq j \leq \frac{t-3}{2} - 1$



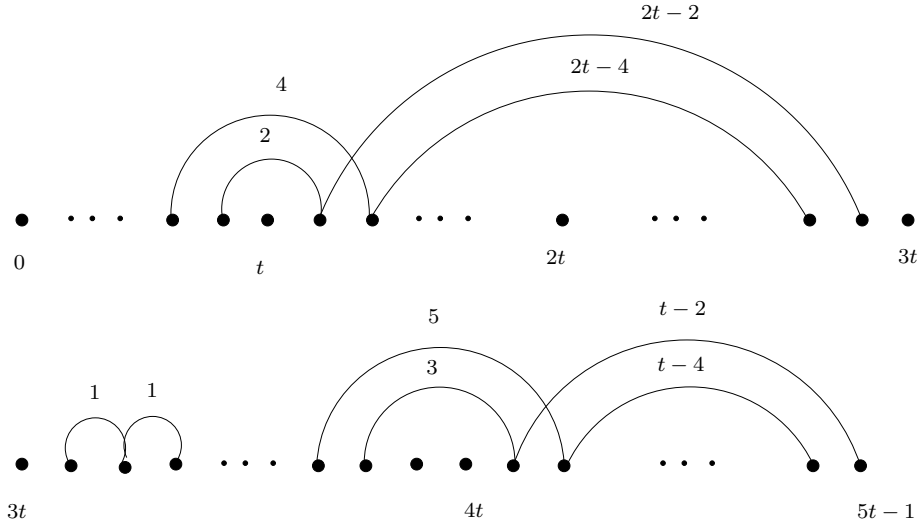
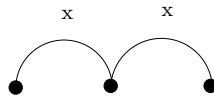


Figure 5.4: Bridges for Case 2

This produces  $(t + \frac{t-3}{2})$   $\delta$  base blocks and  $v \geq 5t - 5$ , which is small enough to overlap with the recursion when  $t \geq 3$ . As shown in figure 5.4 the construction of the odd  $\delta$  blocks requires  $t \geq 6$ . Modifications on this construction to cover  $t = 3$  and 5 will be described following Case 3.

**Case 3:**  $t$  is even

The even pure difference construction of  $\delta$  blocks as described in Case 2 produces  $(t-1)$   $\delta$  base blocks. When  $t$  is even, there are an odd number of  $\delta$  base blocks produced and the bridge when  $i = \frac{t-2}{2}$  will look like



where  $x = t$ . Since the pure difference  $t$  has already been used, this  $\delta$  block may not be constructed.

$\delta$  base blocks:

$$((0, 1), (2t-1-i, 2), (2t-2-2i, 1), (2t+1+i, 2)) \text{ where } 0 \leq i < \frac{t-2}{2} \text{ and } \frac{t-2}{2} < i \leq t-2$$

$$((0, 1), (3t+2, 2), (1, 1), (3t+3, 2)),$$

$((0, 1), (4t + 2 + j, 2), (3 + 2j, 1), (5t - 1 - j, 2))$  where  $0 \leq j \leq \frac{t-4}{2} - 1$

This construction of the  $\delta$  blocks requires  $t \geq 6$ . However,  $(t - 1 + \frac{t-4}{2})$   $\delta$  base blocks are produced which will allow  $v$  to overlap with the recursion when  $t > 8$ . Modifications on this construction and special cases of the recursion to cover  $t = 2, 4, 6, 8$  will be described.

Here are the details for  $t = 2, 3, 4, 5, 6$ , and  $8$ .

**(20,  $v$ ):**

For  $20 \leq v \leq 38$  use Lemma 5.2 and  $(10, h)$  where  $0 \leq h \leq 18$ . When  $v = 18$ , Lemma 3.1 applies. When  $v = 13$ , Lemma 4.3 applies.

$v = 10, 11, 12, 14, 15, 16, 17$  and  $19$ :

Use Lemma 5.3 where  $(a, b, c) = (10, 10, 10)$  and  $v = 0, 1, 2, 4, 5, 6, 7$  and  $9$  respectively.

$v = 9$ :

Use the Recursion with  $x = 8$ .

$4 \leq v \leq 8$ :

Use the Recursion with  $x = 10$ .

**(30,  $v$ ):**

Modify the above construction using  $\delta$  base blocks  $((0, 1), (4, 2), (2, 1), (8, 2))$  and  $((0, 1), (5, 2), (4, 1), (7, 2))$  to cover values  $20 \leq v \leq 58$ .

$v = 19$ :

Use the Recursion with  $x = 13$ .

$v = 18$ : Lemma 3.1 applies.

$v = 16$  and  $17$ :

Use the Recursion with  $x = 14$  and  $12$  respectively.

$11 \leq v \leq 15$ :

Use the Recursion with  $x = 10$ .

$v \leq 10$ : The Recursion applies.

**(40,  $v$ ):**

For  $20 \leq v \leq 78$  modify the above construction using  $\delta$  base blocks:

$((0, 1), (2, 2), (1, 1), (3, 2)), ((0, 1), (0, 2), (3, 1), (14, 2)),$

$((0, 1), (10, 2), (5, 1), (13, 2)), ((0, 1), (17, 2), (2, 1), (19, 2)).$

$v = 19$ : Use the Recursion with  $x = 13$ .

$v \leq 18$ : The Recursion applies.

**(50,  $v$ ):**

For  $20 \leq v \leq 98$  use the above construction replacing the  $\delta$  base blocks with:

$((0, 1), (6, 2), (2, 1), (14, 2)), ((0, 1), (7, 2), (4, 1), (13, 2)), ((0, 1), (8, 2), (6, 1), (12, 2)), ((0, 1), (9, 2), (8, 1), (11, 2)), ((0, 1), (18, 2), (1, 1), (21, 2)), ((0, 1), (22, 2), (3, 1), (23, 2)).$

$v \leq 20$ : The Recursion applies.

**(60,  $v$ ):**

For  $30 \leq v \leq 118$  use the above construction with the additional  $\delta$  base block:

$((0, 1), (15, 2), (12, 1), (27, 2)).$

$v < 30$ : The Recursion applies.

**(80,  $v$ ):**

For  $50 \leq v \leq 158$ : Lemma 5.5 applies.

$v = 49, 48$ : Use the Recursion with  $x = 30$ .

$v \leq 47$ : The Recursion applies. □

## Chapter 6

### Conclusion

**Theorem 6.1** *There exists a  $K_4 - e$  design on  $K_d + v$  when  $d$  is even if and only if*

1.  $5|d(d + 2v - 1)$
2.  $v \leq 2(d - 1)$  and  $v \neq 2d - 3, v \neq 2d - 4$

**Proof** This is a direct result from Lemmas 3.1, 4.1, 4.2, 4.3, 5.1 and 5.5. □

To solve the case when  $d$  is odd, new techniques will need to be applied. The use of difference methods in the construction of  $\delta$  blocks and the use of 1-factors is not applicable when  $d$  is even. However, Lemma 2.1, the Recursion, and Lemma 5.3 are true regardless of the parity of  $d$  and may be useful tools in the future.

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