

Gregarious Path Decompositions of Some Graphs

by

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Abstract

Let G be a simple graph and $f(v)$ a positive integer for each vertex v of G . Form G^f by replacing each v by a set $F(v)$ of $f(v)$ vertices, and each edge uv by complete bipartite graph on bipartition $(F(u), F(v))$. Can we partition G^f into paths of length 2 which are gregarious, that is, meet three different $F(u)$'s?

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List of Abbreviations

PBBG Parity Balanced Bipartite Graph

BGwFD Bipartite Graph with Four Degrees

$L(G)$ Line Graph of G

GDDs Group Divisible Designs

NASCs Necessary and Sufficient Conditions

TFAE The following are equivalent

Chapter 1

Introduction

Let $G = (V, E)$ be a simple graph and $f : V \rightarrow \mathbb{N}$, where $f(v)$ is a positive integer for each vertex v of G . Form the graph G^f by replacing each v by a set $F(v)$ of $f(v)$ vertices, and each edge uv by a complete bipartite graph on bipartition $(F(u), F(v))$. Our question is: “Can we partition G^f into paths of length 2, P_3 , which are *gregarious*, that is, each vertex of P_3 is in a different $F(u)$?”

Example 1.1. Let $G = (V, E)$ be a graph where $|V| = 5$ and $f : V(G) \rightarrow \mathbb{P}$ such that $f : (v_1, v_2, v_3, v_4) \rightarrow (2, 3, 2, 2)$.

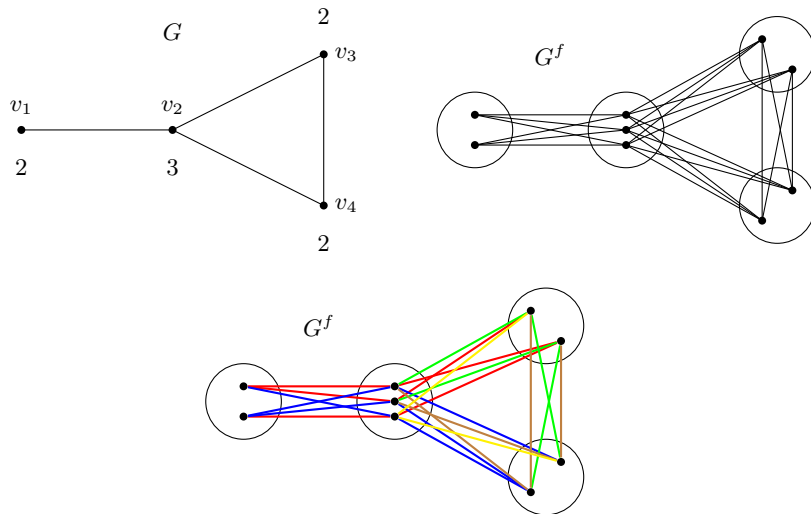


Figure 1.1: Example of Gregarious Path Decomposition

Each color represents a different class of P_3 's.

1.1 History of the problem

Non-Gregarious Case:

In a recent paper [3], the complete solution is given for a decomposition of any complete multipartite graph into paths of lengths 3 and 4. Dr. Hoffman & Dr. Billington introduce the problem “what if we try to solve the same problem with **gregarious** paths?”

Previous work in Gregarious Decompositions:

1. In [5], Dean G. Hoffman & Elizabeth Billington give necessary and sufficient conditions to decompose a complete tripartite graph into gregarious 4-cycles. They use the notion of gregarious decompositions as “a cycle is said to be gregarious if its vertices occur in as many different parts of the multipartite graph as possible”.
2. In [6], Dean G. Hoffman & Elizabeth Billington give the necessary and sufficient conditions for gregarious 4-cycle decompositions of the complete equipartite graph $K_{n(m)}$ (with $n \geq 4$ parts of size m) whenever a 4-cycle decomposition (gregarious or not) is possible, and also of a complete multipartite graph in which all parts but one have the same size.
3. In [7], Benjamin R. Smith give necessary and sufficient conditions for the existence of a gregarious 5-cycle decomposition of the complete equipartite graph $K_{m(n)}$.
4. In [8], Elizabeth J. Billington, Benjamin R. Smith & D.G. Hoffman give necessary and sufficient conditions for gregarious cycle decomposition of the complete equipartite graph $K_{n(m)}$ (with n parts, $n \geq 6$ or $n \geq 8$, of size m) into both 6-cycles and 8-cycles.
5. In [9] Jung R. Cho & Ronald J. Gould give necessary and sufficient conditions for the existence of decompositions of the complete multipartite graph $K_{n(2t)}$ into gregarious 6-cycles if $n \equiv 0, 1, 3$ or $4 \pmod{6}$. They used the method of a complete set of differences in \mathbb{Z}_n .

6. In [10], Jung R. Cho gives another proof of the problem of decomposing the complete multipartite graph $K_{n(2t)}$ into gregarious 6-cycles for the case of $n \equiv 0$ or $3 \pmod{6}$.
7. In [11], Benjamin R. Smith gives necessary and sufficient conditions for the existence of gregarious k - cycle decomposition of a complete equipartite graph, having n parts of size m , and either $n \equiv 0, 1 \pmod{k}$, or k is odd and $m \equiv 0 \pmod{k}$.
8. In [12], Elizabeth J. Billington , Dean G. Hoffman & Chris A. Rodger give necessary and sufficient conditions for decomposing a complete equipartite graph $K_{n(m)}$ with n parts of size m into n -cycles in such a way that each cycle meets each part of $K_{n(m)}$; that is, each cycle is said to be gregarious. Furthermore, they give gregarious decompositions which are also resolvable.
9. In [13], Saad I. El-Zanati, Narong Punnim & Chris A. Rodger give necessary and sufficient conditions for the existence of Gregarious GDDs with Two Associate Classes having block size 3.

Definition 1.2. Let $G = (V, E)$ be simple graph and $h : E \rightarrow \mathbb{P}$. Define $G^{[h]}$ on vertex set V as follows: if $u, v \in V$ and $uv = e \in E$ then put $h(e)$ edges in between u and v in $G^{[h]}$.

Example 1.3. Let's use the example 1.1 and define $h(v_i v_j) := f(v_i) f(v_j)$.

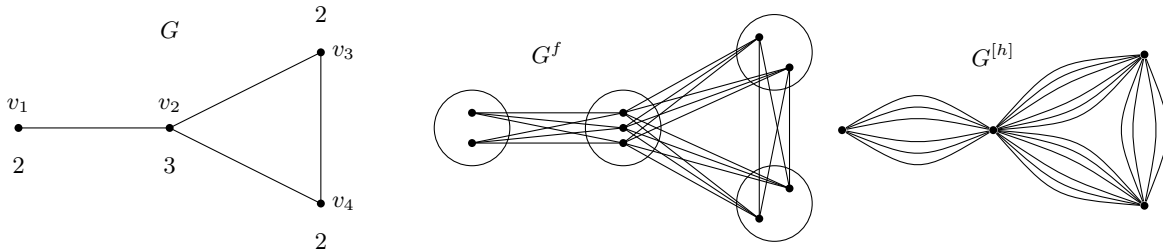


Figure 1.2: Example of G^f and $G^{[h]}$

In $G^{[h]}$, any P_3 will be a gregarious path of length two.

Chapter 2

Tutte's f-Factor Theorem

2.1 Tutte's f -Factor Theorem

Definition 2.1. Let $G = (E, V)$ be a graph. G is called a k -regular graph if for every $v \in V$, $\deg(v) = k$ for some $k \in \mathbb{N}$.

Definition 2.2. Let $G = (V, E)$ be a graph, then

1. A **factor** of a graph G is a spanning subgraph of G .
2. A k -**factor** is a spanning k -regular subgraph.
3. Given a function $f : V(G) \rightarrow \mathbb{Z}$, an f -**factor** of a graph G is a spanning subgraph H such that $d_H(v) = f(v)$ for all $v \in V$.

Example 2.3. Let $G = (V, E)$ be a graph with $f : (v_1, v_2, v_3, v_4, v_5, v_6) \rightarrow (3, 1, 1, 2, 3, 2)$, then we can get the f -factor:

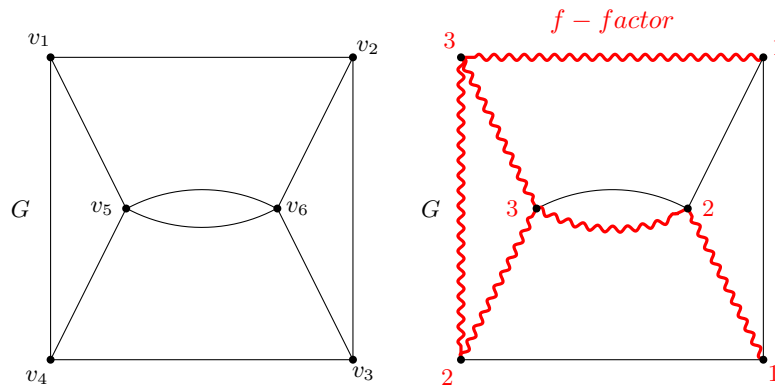


Figure 2.1: Example of an f -factor

Definition 2.4. If $f : V(G) \rightarrow \mathbb{Z}$, define $\bar{f} : V(G) \rightarrow \mathbb{Z}$ by $\bar{f}(v) = \text{deg}_G(v) - f(v)$.

Example 2.5. From the previous example $\bar{f} : (v_1, v_2, v_3, v_4, v_5, v_6) \rightarrow (0, 2, 2, 1, 1, 2)$

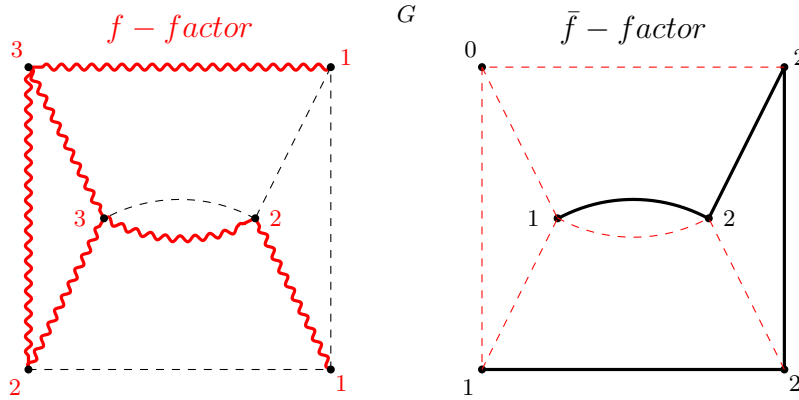


Figure 2.2: Example of an \bar{f} -factor

Definition 2.6. Let $S, T \subseteq V(G)$, with $S \cap T = \emptyset$, then $\lambda(S, T) =$ the set of edges with one end in S and the other end in T .

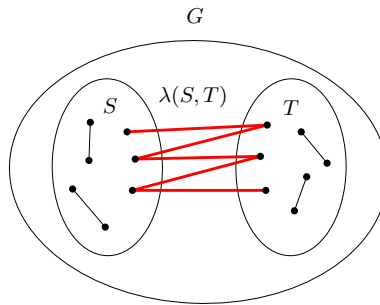


Figure 2.3: Definition of $\lambda(S, T)$

Definition 2.7. $B = (S, T, U)$ is a G -triple if

1. $S, T, U \subseteq V(G)$,
2. $S \cup T \cup U = V(G)$

3. $S \cap T = S \cap U = T \cap U = \emptyset$

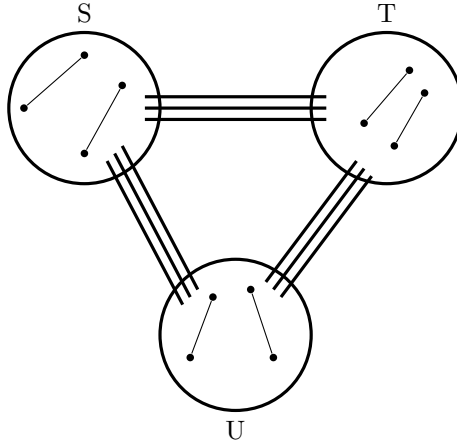


Figure 2.4: Example of a G-triple

Definition 2.8. Let $G = (V, E)$ be a graph and $f : V \rightarrow \mathbb{P}$ be a function.

Then define $f(S) = \sum_{s \in S} f(s)$ for any $S \subseteq V$.

Definition 2.9. By a *component* of B , we mean a component of $G \setminus (S \cup T)$.

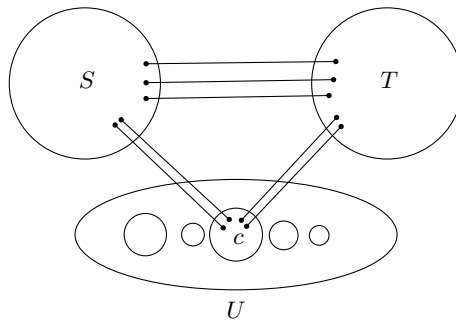


Figure 2.5: Component of B

1. If c is a component of B , let $J(B, f, c) = f(c) + \lambda(c, T)$.
2. c is called **odd** or **even** according to if $J(B, f, c)$ is odd or even.
3. $k(B, f) = \#$ of odd components of B .

Theorem 2.10. [4] G has an f -factor, iff for each G -triple $B = (S, T, U)$

$$k(B, f) + \lambda(S, T) \leq f(S) + \bar{f}(T)$$

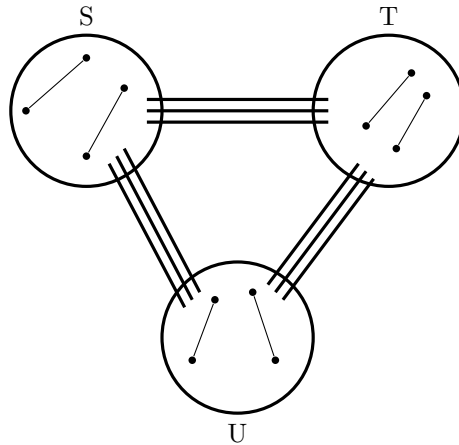


Figure 2.6: Tutte's f-factor Theorem

2.2 Applying Tutte's f-Factor Theorem

Definition 2.11. The *line graph* of a graph G , written $L(G)$, is the graph whose vertices are the edges of G , with $ef \in E(L(G))$ whenever e and f are different edges of G having at least one vertex of G in common.

Let $G = (V, E)$ be a simple graph and $f : V \rightarrow \mathbb{N}$, where $f(v)$ is a positive integer for each vertex v of G . Then, let $h : E \rightarrow \mathbb{P}$ where $h(v_i v_j) := f(v_i) f(v_j)$. In this way we can get G^f and $G^{[h]}$ with given G and f .

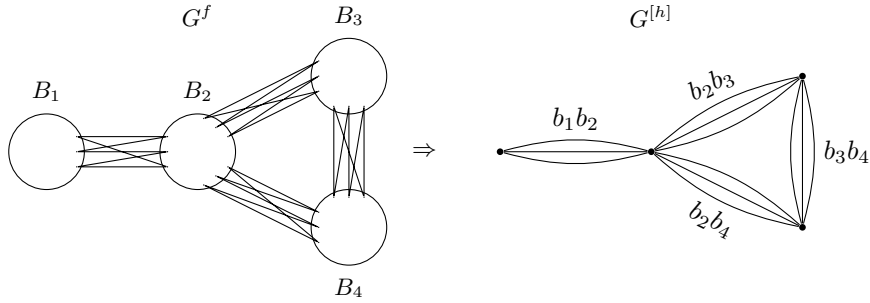


Figure 2.7: G^f and $G^{[h]}$

Now, if we have a gregarious decomposition of G^f into P_3 's (denoted by $G^f \xrightarrow{g} P_3$), then getting decomposition of $G^{[h]}$ into P_3 's (denoted by $G^{[h]} \hookrightarrow P_3$) is trivial.

So $G^f \xrightarrow{g} P_3 \Rightarrow G^{[h]} \hookrightarrow P_3$.

The opposite direction is not true, but it will still give us some of the necessary conditions for $G^f \xrightarrow{g} P_3$. We can apply Tutte's f-factor theorem to solve $G^{[h]} \hookrightarrow P_3$.

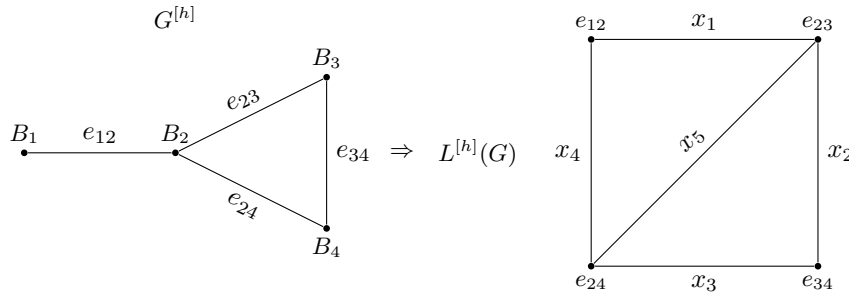


Figure 2.8: $G^{[h]}$ and $L^{[h]}(G)$

In $G^{[h]}$, assume that there are $e_{ij} = b_i b_j$ edges in between any pair of sets of vertices B_i and B_j where $|B_i| = b_i$ for any i . Firstly, get the line graph of G , $L(G)$, then blow the edges of $L(G)$ in such a way that for each vertex $e_{ij} \in V(L(G))$, $deg(e_{ij})$ is $b_i b_j$ in the new graph, $L^{[h]}(G)$. In $L^{[h]}(G)$, each edge will represent a P_3 in $G^{[h]}$. For example, in figure 2.8, x_1 represents the number of P_3 's passing through sets $B_1 \rightarrow B_2 \rightarrow B_3$ in $G^{[h]}$.

Therefore, if we find each x_i , we can find the P_3 decomposition of $G^{[h]}$. To find the x_i 's, we will use Tutte's f-factor theorem. Start with G , get $L(G)$, let M be a sufficiently large number, then get $L^M(G)$ by replacing every edge of $L(G)$ with M edges (fig. 2.9). So if we find an h -factor of $L^M(G)$ such that $h(e_{ij}) = b_i b_j$, then we can get $G^{[h]} \hookrightarrow P_3$.

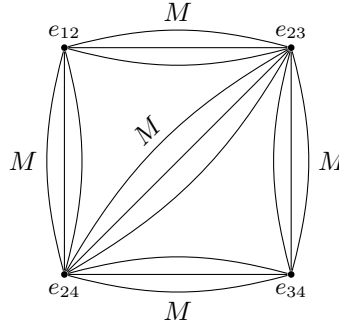


Figure 2.9: $L^M(G)$

Theorem 2.12. $L^M(G)$ has an h -factor for all sufficiently large M , iff for each $L(G)$ -triple $B = (S, T, U)$ where T is independent and $\lambda(T, U) = 0$,

$$k(B, h) + h(T) \leq h(S).$$

Proof. Let M be a sufficiently large number. From Tutte's f-factor theorem, for all $L(G)$ -triples $B = (S, T, U)$:

$$k(B, h) + \lambda_{L^M}(S, T) \leq h(S) + \bar{h}(T)$$

$$k(B, h) + M\lambda_L(S, T) \leq h(S) + M\deg_L(T) - h(T)$$

$$k(B, h) + h(T) - h(S) \leq M(\deg_L(T) - \lambda_L(S, T))$$

Here, $\deg_L(T) \geq \lambda_L(S, T)$ for any $L(G)$ -triple since $\deg_L(T) = \sum_{t \in T} \deg(t)$ and $\lambda_L(S, T) =$ number of edges in between S & T . In addition, when $\deg_L(T) > \lambda_L(S, T)$ the inequality holds since we can choose M sufficiently large and the left hand side of the

inequality doesn't depend on M . So the only problem is when $deg_L(T) = \lambda_L(S, T)$. Which means T is independent and there is no edge in between T & U ($\lambda_L(T, U) = 0$). So the condition we need to check for each $L(G)$ -triple reduces to:

$$k(B, h) + h(T) \leq h(S)$$

when T is independent and $\lambda_L(T, U) = 0$. □

Example 2.13. Let G be the underlying graph in Figure 2.8. Let's find the necessary conditions for $G^{[h]}$ to have a gregarious P_3 decomposition. We need to check

$$k(B, h) + h(T) \leq h(S)$$

for all $L(G)$ triples $B = (S, T, U)$ where T is independent and $\lambda_L(T, U) = 0$.

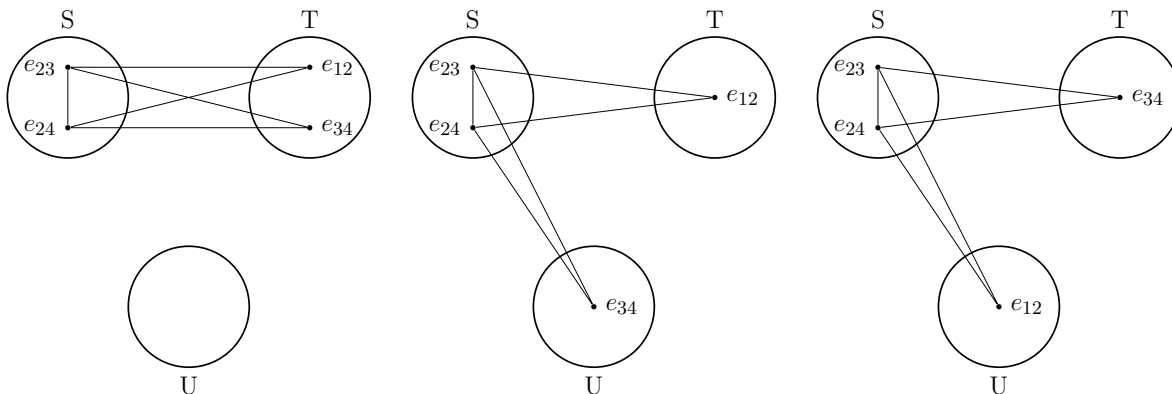


Figure 2.10: Example of how to apply Tutte's f-Factor Theorem

where $deg(e_{ij}) = b_i b_j$. So we can get conditions:

1. $b_1 b_2 + b_3 b_4 \leq b_2 b_3 + b_2 b_4$
2. $b_1 b_2 \leq b_2 b_3 + b_2 b_4$
3. $b_3 b_4 \leq b_2 b_3 + b_2 b_4$

Therefore the necessary condition to decompose $G^{[h]}$ into gregarious P_3 's is:

$$b_1b_2 + b_3b_4 \leq b_2b_3 + b_2b_4$$

since 1 is stronger than 2 and 3.

In summary, if we have a gregarious decomposition of G^f into P_3 's ($G^f \xrightarrow{g} P_3$), then getting a decomposition of $G^{[h]}$ into P_3 's is trivial ($G^{[h]} \hookrightarrow P_3$). So $G^f \xrightarrow{g} P_3 \Rightarrow G^{[h]} \hookrightarrow P_3$.

The opposite direction is not true ($G^{[h]} \hookrightarrow P_3 \not\Rightarrow G^f \xrightarrow{g} P_3$, see the following counter example), but it will still give us some of the necessary conditions for $G^f \xrightarrow{g} P_3$.

Example 2.14. Let $G = S_3$ be a tristar and define $f : V(G) \rightarrow \mathbb{P}$ by

$$f : (a, v_1, v_2, v_3) \rightarrow (2, 1, 1, 1).$$

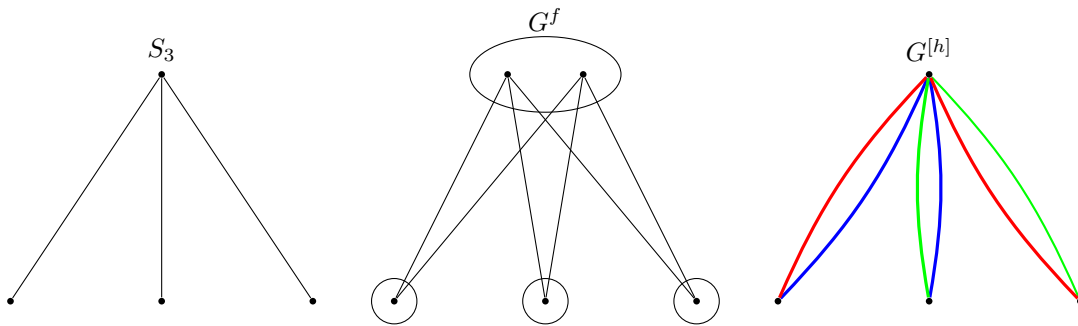


Figure 2.11: Counter Example

G^f doesn't have a gregarious P_3 decomposition since both vertices in the root have odd degree. On the other side, $G^{[h]}$ has a gregarious P_3 decomposition, each color gives a different gregarious P_3 .

Chapter 3

Parity Balanced Bipartite Graphs

Let $a, b \in \mathbb{P}$ and $e \in \mathbb{N}$, and let $\epsilon_a, \epsilon_b \in \{0, 1\}$. We say the simple bipartite graph G on bipartition (A, B) , where $|A| = a$ and $|B| = b$, with e edges, is *parity balanced* with parameters $(a, b, e, \epsilon_a, \epsilon_b)$ if

$\forall u \in A, \deg(u) \equiv \epsilon_a \pmod{2}$, and further $\forall v \in A, |\deg(u) - \deg(v)| \leq 2$,

$\forall u \in B, \deg(u) \equiv \epsilon_b \pmod{2}$, and further $\forall v \in B, |\deg(u) - \deg(v)| \leq 2$.

We will give necessary and sufficient conditions on the parameters $(a, b, e, \epsilon_a, \epsilon_b)$ for the existence of such graphs.

3.1 Introduction

All the graphs are *simple*, i.e., they have no loops or multiple edges. Let \mathbb{P} be the set of positive integers and \mathbb{N} be the set of non-negative integers.

Definition 3.1. The integer vector (x_1, x_2, \dots, x_t) is said to be *balanced* if $|x_i - x_j| \leq 1$ for all $1 \leq i, j \leq t$. Two vectors are *equivalent* if one can be obtained from the other by permuting the entries.

Definition 3.2. Let G be a bipartite graph on bipartition (A, B) . If for all $v \in A, \deg_G(v) = d_1$ and for all $w \in B, \deg_G(w) = e_1$, then we will call G a (d_1, e_1) -*regular* bipartite graph.

The following lemmas are proved in [1], p. 399.

Lemma 3.3. Let v and w be balanced vectors with the same number of coordinates. Then, for some vector w' equivalent to w , $v + w'$ is balanced.

Lemma 3.4. Let $a, b \in \mathbb{P}$, and let $e \leq ab$ be a non-negative integer. Then there is a bipartite graph G on bipartition (A, B) with both $(\deg_G(v) \mid v \in A)$ and $(\deg_G(y) \mid y \in B)$ balanced.

3.2 Bipartite Graphs with Four Degrees

The theorems we will be proving here can be proven by using the Ryser-Gale theorem ([2], p. 185), but the proof is much harder.

Theorem 3.5. *Let $a_1, a_2, b_1, b_2, d_1, d_2, e_1, e_2$ be non-negative integers. Then:*

There is a simple bipartite graph on bipartition (A, B) , where A consists of a_1 vertices of degree d_1 and a_2 vertices of degree d_2 , and B consists of b_1 vertices of degree e_1 and b_2 vertices of degree e_2 , if and only if

$$(*) a_1d_1 + a_2d_2 = b_1e_1 + b_2e_2$$

1. $a_1d_1 \leq a_1b_1 + b_2e_2$, or, equivalently, $b_1e_1 \leq a_1b_1 + a_2d_2$
2. $a_1d_1 \leq a_1b_2 + b_1e_1$, or, equivalently, $b_2e_2 \leq a_1b_2 + a_2d_2$
3. $b_1e_1 \leq a_2b_1 + a_1d_1$, or, equivalently, $a_2d_2 \leq a_2b_1 + b_2e_2$
4. $b_2e_2 \leq a_2b_2 + a_1d_1$, or, equivalently, $a_2d_2 \leq a_2b_2 + b_1e_1$
5. either $a_1 = 0$, or $d_1 \leq b_1 + b_2$
6. either $a_2 = 0$, or $d_2 \leq b_1 + b_2$
7. either $b_1 = 0$, or $e_1 \leq a_1 + a_2$
8. either $b_2 = 0$, or $e_2 \leq a_1 + a_2$

Necessity:

Proof. Each side of $(*)$ counts the total number of edges, hence they must be equal. Conditions 5 - 8 come from the fact that maximum degree of any vertex is less than the number of vertices in the other part. Now for conditions 1 - 4:

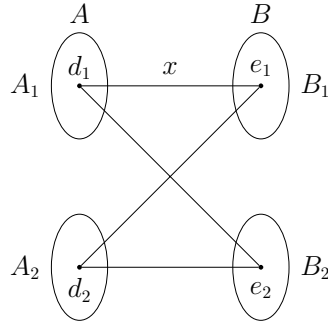


Figure 3.1: Bipartite Graph with Four Degrees

For $i = 1, 2$, let A_i (resp. B_i) be the vertices of degree d_i (resp. e_i) in A_i (resp. B_i). In addition, let's assume that there are x edges in between vertices of A_1 and B_1 . If we look at the other pairs, we will get:

	B_1	B_2
A_1	x	$a_1 d_1 - x$
A_2	$b_1 e_1 - x$	$a_2 d_2 - b_1 e_1 + x$ $=$ $b_2 e_2 - a_1 d_1 + x$

Table 3.1: Distribution of the edges

Using table 3.1, we can get the following inequalities:

$$0 \leq x \leq a_1 b_1$$

$$0 \leq a_1 d_1 - x \leq a_1 b_2$$

$$0 \leq b_1 e_1 - x \leq a_2 b_1$$

$$0 \leq a_2 d_2 - b_1 e_1 + x \leq a_2 b_2$$

So, we can get:

$$\begin{aligned}
0 &\leq x \leq a_1 b_1 \\
a_1 d_1 - a_1 b_2 &\leq x \leq a_1 d_1 \\
b_1 e_1 - a_2 b_1 &\leq x \leq b_1 e_1 \\
b_1 e_1 - a_2 d_2 &\leq x \leq a_2 b_2 - a_2 d_2 + b_1 e_1
\end{aligned}$$

We can get sixteen inequalities on the variables $(a_1, a_2, b_1, b_2, d_1, d_2, e_1, e_2)$ from above since we have x in the middle of all of the four inequalities. If we use the left side of the first inequality and right sides of the all them, then we can get:

$$\begin{aligned}
0 &\leq a_1 b_1 \\
0 &\leq a_1 d_1 \\
0 &\leq b_1 e_1 \\
0 &\leq a_2 b_2 - a_2 d_2 + b_1 e_1 \Rightarrow a_2 d_2 \leq a_2 b_2 + b_1 e_1, \text{ cond. } 4 \checkmark
\end{aligned}$$

From the second one:

$$\begin{aligned}
a_1 d_1 - a_1 b_2 &\leq a_1 b_1 \Rightarrow d_1 \leq b_1 + b_2 \\
a_1 d_1 - a_1 b_2 &\leq a_1 d_1 \Rightarrow 0 \leq a_1 b_1 \\
a_1 d_1 - a_1 b_2 &\leq b_1 e_1 \Rightarrow a_1 d_1 \leq a_1 b_2 + b_1 e_1, \text{ cond. } 2 \checkmark \\
a_1 d_1 - a_1 b_2 &\leq a_2 b_2 - a_2 d_2 + b_1 e_1 \Rightarrow a_1 d_1 + a_2 d_2 \leq a_2 b_2 + b_1 e_1 + a_1 b_2
\end{aligned}$$

If we use (*), we will get $a_1 d_1 + a_2 d_2 = b_1 e_1 + b_2 e_2 \leq a_2 b_2 + b_1 e_1 + a_1 b_2$, so this reduces to $e_1 \leq a_1 + a_2$.

From the third one:

$$\begin{aligned} b_1 e_1 - a_2 b_1 &\leq a_1 b_1 && \Rightarrow e_1 \leq a_1 + a_2 \\ b_1 e_1 - a_2 b_1 &\leq a_1 d_1 && \Rightarrow b_1 e_1 \leq a_2 b_1 + a_1 d_1, \text{ cond. 3 } \checkmark \\ b_1 e_1 - a_2 b_1 &\leq b_1 e_1 && \Rightarrow 0 \leq a_2 b_1 \\ b_1 e_1 - a_2 b_1 &\leq a_2 b_2 - a_2 d_2 + b_1 e_1 && \Rightarrow d_2 \leq b_1 + b_2 \end{aligned}$$

From the fourth one:

$$\begin{aligned} b_1 e_1 - a_2 d_2 &\leq a_1 b_1 && \Rightarrow b_1 e_1 \leq a_1 b_1 + a_2 d_2, \text{ cond. 1 } \checkmark \\ b_1 e_1 - a_2 d_2 &\leq a_1 d_1 && \Rightarrow b_1 e_1 \leq a_1 d_1 + a_2 d_2 \\ b_1 e_1 - a_2 d_2 &\leq b_1 e_1 && \Rightarrow 0 \leq a_2 d_2 \\ b_1 e_1 - a_2 d_2 &\leq a_2 b_2 - a_2 d_2 + b_1 e_1 && \Rightarrow 0 \leq a_2 b_2 \end{aligned}$$

In the second equation, if we use $a_1 d_1 + a_2 d_2 = b_1 e_1 + b_2 e_2$, then we will get

$$b_1 e_1 \leq b_1 e_1 + b_2 e_2 \Rightarrow 0 \leq b_2 e_2.$$

□

Sufficiency:

Proof. Assume there is bipartite graph (A, B) satisfying the necessary conditions and there are x edges in between the vertices of A_1 and B_1 like in figure 3.1. Therefore, we can find the number of edges in between other vertices using the remaining edges like in table 3.1.

Now using the construction in [1] on pg. 399,

1. distribute x edges on a_1 vertices with balanced degrees.
2. distribute $a_1 d_1 - x$ on a_1 vertices with balanced degrees.

So we will get two balanced vectors with the same number of entries.

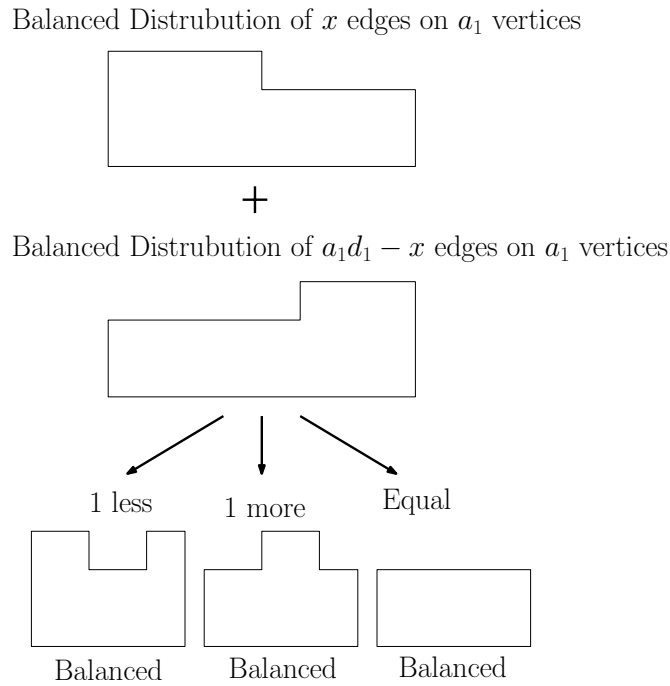


Figure 3.2: Adding balanced distributions

At the end, we will have one of these three cases from figure 3.2 and all of them will still have balanced distributions since both distributions were balanced to begin with. However, the first two cases are impossible, since we have a balanced distribution of $a_1 d_1$ edges on a_1 vertices, this means that each vertex will be incident with d_1 edges. In the same manner we can prove that we can distribute the remaining edges with the desired degrees. □

3.3 Parity balanced Bipartite Graphs

Definition 3.6. Let $a, b \in \mathbb{P}$ and $e \in \mathbb{N}$, and let $\epsilon_a, \epsilon_b \in \{0, 1\}$. We say the bipartite graph G with e edges on bipartition (A, B) , with $|A| = a$ and $|B| = b$, is *parity balanced* with parameters $(a, b, e, \epsilon_a, \epsilon_b)$ if

1. $\forall u \in A, \deg(u) \equiv \epsilon_a \pmod{2}$ and further $\forall v \in A, |\deg(u) - \deg(v)| \leq 2$.
2. $\forall u \in B, \deg(u) \equiv \epsilon_b \pmod{2}$ and further $\forall v \in B, |\deg(u) - \deg(v)| \leq 2$.

Example 3.7. Let $|A| = a = 6, |B| = b = 5, e = 14, \epsilon_a = 1$ and $\epsilon_b = 0$.

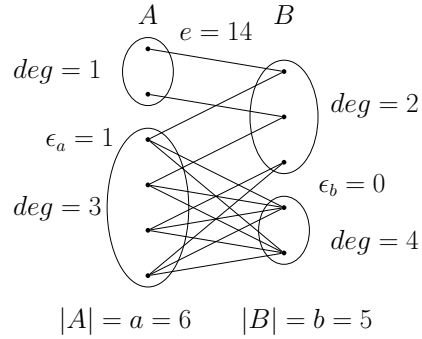


Figure 3.3: Example of a parity balanced bipartite graph

Definition 3.8. If A and B are disjoint sets, we denote $K_{A,B}$ to be the complete bipartite graph on bipartition (A, B) .

Definition 3.9. Let $K_{A,B}$ be a complete bipartite graph on bipartition (A, B) . A *bipartite complement* of a bipartite graph G on bipartition (A, B) with edge set E is the bipartite graph G' on bipartition (A, B) with the edge set E' where $E' = E(K_{A,B}) \setminus E$.

Fact 3.10. If G is a parity balanced bipartite graph with parameters $(a, b, e, \epsilon_a, \epsilon_b)$, then G' is a parity balanced bipartite graph with parameters $(a, b, e' = ab - e, \epsilon'_a, \epsilon'_b)$ where $\epsilon_a + \epsilon'_a \equiv b \pmod{2}$ and $\epsilon_b + \epsilon'_b \equiv a \pmod{2}$

Example 3.11. The bipartite complement of G is G' with parameters $(a = 6, b = 5, e' = ab - 14 = 30 - 14 = 16, \epsilon'_a = 0, \epsilon'_b = 0)$:

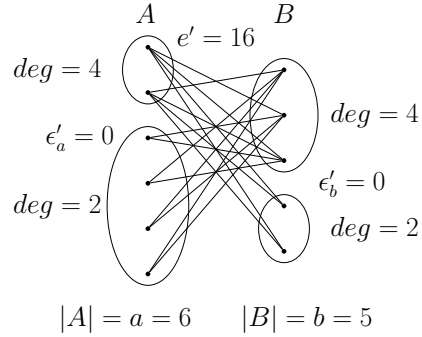


Figure 3.4: Example of a bipartite complement

where $\epsilon_a + \epsilon'_a = 1 + 0 = 1 \equiv 5 \pmod{2}$ and $\epsilon_b + \epsilon'_b = 0 + 0 \equiv 6 \pmod{2}$.

Theorem 3.12. Let $a, b \in \mathbb{P}$, $e \in \mathbb{N}$, $\epsilon_a, \epsilon_b, \epsilon'_a, \epsilon'_b \in \{0, 1\}$,

$\epsilon_a + \epsilon'_a \equiv b \pmod{2}$, $\epsilon_b + \epsilon'_b \equiv a \pmod{2}$.

Then, there is a parity balanced bipartite graph G on bipartition (A, B) with parameters $(a, b, e, \epsilon_a, \epsilon_b)$ if and only if

$\epsilon_a a \leq e \leq ab - \epsilon'_a a$, $\epsilon_b b \leq e \leq ab - \epsilon'_b b$, and all of these are congruent $\pmod{2}$, with the following exceptions:

e	a	b	ϵ_a	ϵ'_a	ϵ_b	ϵ'_b
2	2	2	0	0	0	0
	2	3	0	1	0	0
	3	2	0	0	0	1
	odd ≥ 3	odd ≥ 3	0	1	0	1
	odd ≥ 3	even ≥ 3	0	0	0	1
			0	0	1	0
	even ≥ 3	odd ≥ 3	0	1	0	0
			1	0	0	0
even ≥ 3	even ≥ 3	1	1	1	1	
		0	0	0	0	
$ab-2$	2	2	0	0	0	0
	2	3	1	0	0	0
	3	2	0	0	1	0
	odd ≥ 3	odd ≥ 3	1	0	1	0
	odd ≥ 3	even ≥ 3	0	0	1	0
			0	0	0	1
	even ≥ 3	odd ≥ 3	1	0	0	0
			0	1	0	0
even ≥ 3	even ≥ 3	1	1	1	1	
		0	0	0	0	

Table 3.2: Exceptions for Theorem [3.12](#)

Necessity:

Proof. For any $u \in A$ we have $\deg_G(u) + \deg_{G'}(u) = b$ so $\deg_G(u) + \deg_{G'}(u) \equiv b \pmod{2}$ where $\deg_G(u) \equiv \epsilon_a \pmod{2}$ and $\deg_{G'}(u) \equiv \epsilon'_a \pmod{2}$ by definition, and $\epsilon_a + \epsilon'_a \equiv b \pmod{2}$ follows. In the same way, we can get $\epsilon_b + \epsilon'_b \equiv a \pmod{2}$.

To get $\epsilon_a a \leq e \leq ab - \epsilon'_a a$ and $\epsilon_b b \leq e \leq ab - \epsilon'_b b$, if one of $\epsilon_a, \epsilon'_a, \epsilon_b, \epsilon'_b$ is 1, then we have to have enough edges in either G or G' .

Finally, to get $\epsilon_a a, e, ab - \epsilon'_a a, \epsilon_b b, ab - \epsilon'_b b$ all congruent $\pmod{2}$;

$$e = \sum_{u \in A} \deg_G(u) \equiv a\epsilon_a \pmod{2}, \text{ and}$$

$$\epsilon_a + \epsilon'_a \equiv b \pmod{2} \Rightarrow a\epsilon_a + a\epsilon'_a \equiv ab \pmod{2} \Rightarrow a\epsilon_a \equiv ab - a\epsilon'_a \pmod{2}.$$

In the same way we can get the other conditions. For the exceptions, it is easy to prove that there is no parity balanced bipartite graph with parameters given in table 4.1. Figure 3.6 shows all possible parity balanced bipartite graphs with 2 edges and the ones with $ab - 2$ edges will be bipartite complement of these graphs. \square

Sufficiency:

Proof. We can use theorem 3.5 for this proof. Define $n, m, q_a, r_a, q_b, r_b \in \mathbb{N}$ by,

$$e = 2n + \epsilon_a a = 2m + \epsilon_b b$$

$$n = aq_a + r_a, m = bq_b + r_b$$

$$0 \leq r_a \leq a - 1, 0 \leq r_b \leq b - 1.$$

So $e = 2aq_a + 2r_a + \epsilon_a a = 2bq_b + 2r_b + \epsilon_b b$.

$$q_a = \frac{e - \epsilon_a a - 2r_a}{2a} = \frac{e - \epsilon_a a}{2a} - \frac{r_a}{a} = \left\lfloor \frac{e - \epsilon_a a}{2a} \right\rfloor. \text{ In the same way, } q_b = \left\lfloor \frac{e - \epsilon_b b}{2b} \right\rfloor.$$

Now let's translate this problem into "bipartite graphs with four degrees" since we already know NASCs for those graphs (figure 3.5).

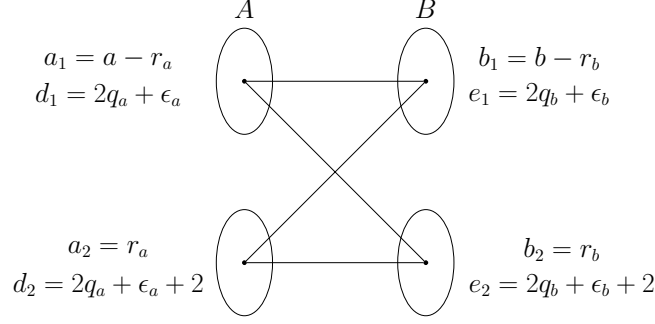


Figure 3.5: Translating PBBG to a BGwFDs

Note that (*) holds since:

$$(a - r_a)(2q_a + \epsilon_a) + r_a(2q_a + \epsilon_a) = (b - r_b)(2q_b + \epsilon_b) + r_b(2q_b + \epsilon_b)$$

Case 1: Assume $a_2 = 0 = b_2$, then we will get a (d_1, e_1) -regular bipartite graph. Since $a_2 = 0$, and $b_2 = 0$, we only need to prove 1, 5 and 7 in theorem 3.5. Let's start proving conditions 5 and 7 which say:

$$\begin{aligned} d_1 &\leq b_1 + b_2 \\ 2q_a + \epsilon_a &\leq b - r_b + r_b = b \end{aligned}$$

and

$$\begin{aligned} e_1 &\leq a_1 + a_2 \\ 2q_b + \epsilon_b &\leq a - r_a + r_a = a \end{aligned}$$

So for 5 if we prove $2q_a + \epsilon_a \leq b$, we are done. Using $2q_a = \frac{e - \epsilon_a a - 2r_a}{a} = \frac{e}{a} - \epsilon_a - \frac{2r_a}{a}$; we can get $2q_a + \epsilon_a = \frac{e}{a} - \epsilon_a - \frac{2r_a}{a} + \epsilon_a = \frac{e}{a} - \frac{2r_a}{a} \leq b - \frac{2r_a}{a} < b$. We can prove 7 in the same way.

Now, let's prove 1:

$$a_1 d_1 \leq a_1 b_1 + b_2 e_2$$

$$a_1 d_1 \leq a_1 b_1 \text{ since } b_2 = 0$$

$$d_1 \leq b_1$$

$$2q_a + \epsilon_a \leq b \text{ and we just proved this in 5}$$

Case 2: Assume $a_2 = 0$ and $b_2 \neq 0$ ($a_2 \neq 0$ and $b_2 = 0$ is just the symmetric case). Since $a_2 = 0$, we only need to prove 1, 2, 5, 7 and 8 in theorem 3.5. 5 and 7 are the same as in Case 1. 1 and 2 will reduce to 7 and 8, respectively since $a_2 = 0$. So just proving 8 is enough which says:

$$e_2 \leq a_1 + a_2$$

$$2q_b + \epsilon_b + 2 \leq a_1 = a \text{ since } a_2 = 0$$

So $2q_b = \frac{e}{b} - \epsilon_b - \frac{2r_b}{b}$, using this:

$$2q_b + \epsilon_b + 2 = \frac{e}{b} - \epsilon_b - \frac{2r_b}{b} + \epsilon_b + 2 = \frac{e}{b} - \frac{2r_b}{b} + 2 \leq a - \epsilon'_b - \frac{2r_b}{b} + 2 < a + 2.$$

The only problem is when $a = 2q_b + \epsilon_b + 1$, then $\epsilon_b + \epsilon'_b \equiv a = 2q_b + \epsilon_b + 1 \pmod{2}$.

So $\epsilon'_b = 1$. In this case:

$$2q_b + \epsilon_b + 2 = \frac{e}{b} - \epsilon_b - \frac{2r_b}{b} + \epsilon_b + 2 = \frac{e}{b} - \frac{2r_b}{b} + 2 \leq a - \epsilon'_b - \frac{2r_b}{b} + 2 < a + 1.$$

Case 3: We can assume $a_2 \neq 0 \neq b_2$. Let's start proving conditions 5 through 8 in theorem 3.5. So 5 and 6 say:

$$d_1 \leq b_1 + b_2$$

$$2q_a + \epsilon_a \leq b - r_b + r_b = b$$

and

$$\begin{aligned} d_2 &\leq b_1 + b_2 \\ 2q_a + \epsilon_a + 2 &\leq b - r_b + r_b = b \end{aligned}$$

If we prove $2q_a + \epsilon_a + 2 \leq b$, this will cover both cases. However, this is the same as (8) in Case 2, just switch a and b . We can prove (7) and (8) in the same way.

Now let's prove conditions 1 through 4 of theorem 3.5.

For 1, we need to prove:

$$\begin{aligned} a_1 d_1 &\leq a_1 b_1 + b_2 e_2 \\ (a - r_a)(2q_a + \epsilon_a) &\leq (a - r_a)(b - r_b) + r_b(2q_b + \epsilon_b + 2) \end{aligned}$$

First, suppose $r_b \leq b - 2q_a - \epsilon_a$, then we get:

$$\begin{aligned} (a - r_a)(2q_a + \epsilon_a) &\leq (a - r_a)(b - r_b) + r_b(2q_b + \epsilon_b + 2) \\ (a - r_a)(2q_a + \epsilon_a - b + r_b) &\leq r_b(2q_b + \epsilon_b + 2) \\ (a - r_a)(r_b - (b - 2q_a - \epsilon_a)) &\leq r_b(2q_b + \epsilon_b + 2) \end{aligned}$$

So $r_b - (b - 2q_a - \epsilon_a) \leq 0$ and the inequality is automatically satisfied since $a - r_a > 0$, $r_b > 0$, $2q_b + \epsilon_b + 2 > 0$.

So we can assume $r_b \geq b - 2q_a - \epsilon_a + 1$. In addition, recall that $e = 2n + \epsilon_a a = 2aq_a + 2r_a + \epsilon_a a$.

Need to prove:

$$\begin{aligned}
a_1 d_1 &\leq a_1 b_1 + b_2 e_2 \\
(a - r_a)(2q_a + \epsilon_a) &\leq (a - r_a)(b - r_b) + r_b(2q_b + \epsilon_b + 2) \\
2aq_a + a\epsilon_a - r_a(2q_a + \epsilon_a) &\leq ab - ar_b - br_a + r_a r_b + r_b(2q_b + \epsilon_b + 2) \\
2aq_a + a\epsilon_a - r_a(2q_a + \epsilon_a) + 2r_a - 2r_a &\leq ab - ar_b - br_a + r_a r_b + r_b(2q_b + \epsilon_b + 2) \\
(2aq_a + a\epsilon_a + 2r_a) - r_a(2q_a + \epsilon_a + 2) &\leq ab - ar_b - br_a + r_a r_b + r_b(2q_b + \epsilon_b + 2) \\
e + r_a(b - (2q_a + \epsilon_a + 2)) + r_b(a - (2q_b + \epsilon_b + 2)) - r_a r_b &\leq ab
\end{aligned}$$

So if we show $e + r_a(b - (2q_a + \epsilon_a + 2)) + r_b(a - (2q_b + \epsilon_b + 2)) - r_a r_b \leq ab$, we are done.

Using $r_b \geq b - 2q_a - \epsilon_a + 1$, we can get $r_a(b - (2q_a + \epsilon_a + 2)) < r_a r_b$. So

$$\begin{aligned}
e + r_a(b - (2q_a + \epsilon_a + 2)) + r_b(a - (2q_b + \epsilon_b + 2)) - r_a r_b &< e + r_a r_b + r_b(a - (2q_b + \epsilon_b + 2)) - r_a r_b \\
&= e + r_b(a - (2q_b + \epsilon_b + 2))
\end{aligned}$$

where we can use, $e = (b - r_b)(2q_b + \epsilon_b) + r_b(2q_b + \epsilon_b + 2)$.

$$\begin{aligned}
&= e + r_b(a - (2q_b + \epsilon_b + 2)) \\
&= (b - r_b)(2q_b + \epsilon_b) + r_b(2q_b + \epsilon_b + 2) + r_b(a - (2q_b + \epsilon_b + 2)) \\
&= (b - r_b)(2q_b + \epsilon_b) + ar_b
\end{aligned}$$

here using the fact that $2q_b + \epsilon_b \leq a$ (which we just proved in 7), we will get:

$$\begin{aligned}
&= (b - r_b)(2q_b + \epsilon_b) + ar_b \\
&< (b - r_b)a + ar_b = ab.
\end{aligned}$$

Now let's prove 2:

$$\begin{aligned} a_1 d_1 &\leq a_1 b_2 + b_1 e_1 \\ (a - r_a)(2q_a + \epsilon_a) &\leq (a - r_a)r_b + (b - r_b)(2q_b + \epsilon_b) \end{aligned}$$

First suppose $2q_a + \epsilon_a \leq r_b$, then we get:

$$\begin{aligned} (a - r_a)(2q_a + \epsilon_a) &\leq (a - r_a)r_b + (b - r_b)(2q_b + \epsilon_b) \\ (a - r_a)(2q_a + \epsilon_a - r_b) &\leq (b - r_b)(2q_b + \epsilon_b). \end{aligned}$$

So $2q_a + \epsilon_a - r_b \leq 0$ and the inequality is automatically satisfied since $a - r_a > 0$, $(b - r_b) > 0$, $2q_b + \epsilon_b \geq 0$.

We can assume $2q_a + \epsilon_a + 1 \geq r_b$. We need to prove:

$$\begin{aligned} a_1 d_1 &\leq a_1 b_2 + b_1 e_1 \\ (a - r_a)(2q_a + \epsilon_a) &\leq (a - r_a)r_b + (b - r_b)(2q_b + \epsilon_b) \\ 2aq_a + a\epsilon_a - r_a(2q_a + \epsilon_a) &\leq ar_b - r_a r_b + 2bq_b + b\epsilon_b - r_b(2q_b + \epsilon_b) \\ 2aq_a + a\epsilon_a - r_a(2q_a + \epsilon_a) + 2r_a - 2r_a &\leq ar_b - r_a r_b + 2bq_b + b\epsilon_b - r_b(2q_b + \epsilon_b) + 2r_b - 2r_b \\ (2aq_a + a\epsilon_a + 2r_a) - r_a(2q_a + \epsilon_a + 2) &\leq ar_b - r_a r_b + (2bq_b + b\epsilon_b + 2r_b) - r_b(2q_b + \epsilon_b + 2) \\ e - r_a(2q_a + \epsilon_a + 2) &\leq ar_b - r_a r_b + e - r_b(2q_b + \epsilon_b + 2) \\ r_a r_b + r_b(2q_b + \epsilon_b + 2) &\leq ar_b + r_a(2q_a + \epsilon_a + 2) \end{aligned}$$

If we show $r_a r_b + r_b(2q_b + \epsilon_b + 2) \leq ar_b + r_a(2q_a + \epsilon_a + 2)$, then we have shown (2). Using $2q_a + \epsilon_a + 1 \geq r_b$ we can get $r_a r_b < r_a(2q_a + \epsilon_a + 2)$. In addition, we can use the previously

proved fact in 8 that $2q_b + \epsilon_b + 2 \leq a$. So

$$\begin{aligned}
r_a r_b + r_b(2q_b + \epsilon_b + 2) &\leq r_a(2q_a + \epsilon_a + 2) + r_b(2q_b + \epsilon_b + 2) \\
&\leq r_a(2q_a + \epsilon_a + 2) + r_b a \\
&= ar_b + r_a(2q_a + \epsilon_a + 2)
\end{aligned}$$

The proof of 3 is exactly the same as the proof of 2, if we switch parts $A \longleftrightarrow B$.

Now, let's prove the last condition, 4.

We need to prove $b_2 e_2 \leq a_2 b_2 + a_1 d_1$, or, equivalently, $a_2 d_2 \leq a_2 b_2 + b_1 e_1$.

$$\begin{aligned}
b_2 e_2 &\leq a_2 b_2 + a_1 d_1 \\
r_b(2q_b + \epsilon_b + 2) &\leq r_a r_b + (a - r_a)(2q_a + \epsilon_a) \\
r_b(2q_b + \epsilon_b + 2 - r_a) &\leq (a - r_a)(2q_a + \epsilon_a)
\end{aligned}$$

which is equivalent to

$$\begin{aligned}
a_2 d_2 &\leq a_2 b_2 + b_1 e_1 \\
r_a(2q_a + \epsilon_a + 2) &\leq r_a r_b + (b - r_b)(2q_b + \epsilon_b) \\
r_a(2q_a + \epsilon_a + 2 - r_b) &\leq (b - r_b)(2q_b + \epsilon_b).
\end{aligned}$$

First, assume $2q_b + \epsilon_b + 2 \leq r_a$ or $2q_a + \epsilon_a + 2 \leq r_b$.

Then $b_2 e_2 \leq a_2 b_2 + a_1 d_1$ or $a_2 d_2 \leq a_2 b_2 + b_1 e_1$ will be automatically satisfied.

So we can assume $2q_b + \epsilon_b + 2 > r_a$ and $2q_a + \epsilon_a + 2 > r_b \Rightarrow 2q_a + \epsilon_a + 1 \geq r_b$.

If we turn back to the problem and use the fact, which follows from 8, that $2q_b + \epsilon_b + 2 \leq a$,

then:

$$\begin{aligned}
r_b(2q_b + \epsilon_b + 2) &\leq r_b a \\
&= r_a r_b + (a - r_a) r_b \\
&\leq r_a r_b + (a - r_a)(2q_a + \epsilon_a + 1)
\end{aligned}$$

So the only problem is when $r_b = 2q_a + \epsilon_a + 1$. Similarly, we can assume $r_a = 2q_b + \epsilon_b + 1$.

Need to show:

$$\begin{aligned}
b_2 e_2 &\leq a_2 b_2 + a_1 d_1 \\
r_b(2q_b + \epsilon_b + 2) &\leq r_a r_b + (a - r_a)(2q_a + \epsilon_a) \\
(2q_a + \epsilon_a + 1)(2q_b + \epsilon_b + 2) &\leq (2q_a + \epsilon_a + 1)(2q_b + \epsilon_b + 1) + (a - r_a)(2q_a + \epsilon_a) \\
2q_a + \epsilon_a + 1 &\leq (a - r_a)(2q_a + \epsilon_a) \\
1 &\leq (a - r_a - 1)(2q_a + \epsilon_a) \\
1 &\leq (a - r_a - 1)(r_b - 1)
\end{aligned}$$

Here $r_b \geq 1$ since $r_a \neq 0 \neq r_b$. In this case both are positive and we can assume $a - r_a \geq 1$ since $0 \leq r_a < a$. Therefore, we only need to prove $r_b \neq 1$ or $a - r_a \neq 1$. First, suppose $r_b = 1$, then $r_b = 2q_a + \epsilon_a + 1 = 1$ so $q_a = 0$ and $\epsilon_a = 0$.

$$\begin{aligned}
e = 2aq_a + 2ra + a\epsilon_a &= 2bq_b + 2r_b + b\epsilon_b \\
2r_a &= 2bq_b + 2 + b\epsilon_b \\
2(2q_b + \epsilon_b + 1) &= b(2q_b + \epsilon_b) + 2 \\
2(2q_b + \epsilon_b) &= b(2q_b + \epsilon_b) \\
0 &= (b - 2)(2q_b + \epsilon_b) \\
0 &= (b - 2)(r_a - 1)
\end{aligned}$$

Therefore, in this case, either $b = 2$ or $r_a = 1$. If $r_a = 1$, then $e = 2r_a = 2$ and this is not possible since when $e = 2$ there are only two bipartite graphs with 2 edges (see figure 3.6) and both of them have either $b_2 = 0$ or $a_2 = 0 = b_2$, which contradicts the assumption $a_2 \neq 0 \neq b_2$.

We can get the exceptions in table 4.1 with parameters $(a, b, e = 2, \epsilon_a, \epsilon'_a, \epsilon_b, \epsilon'_b)$ easily since no other bipartite graphs exist with 2 edges but the ones in figure 3.6.

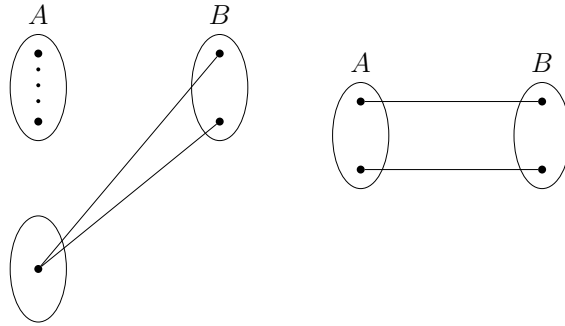


Figure 3.6: Bipartite graphs with 2 edges

Now, assume $b = 2$ where $r_b = 1$ and $r_a \geq 2$.

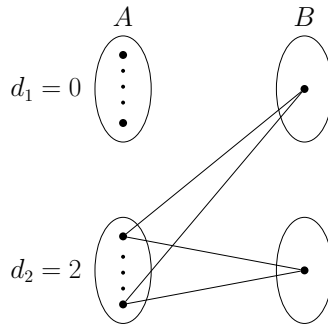


Figure 3.7: Exception for $b = 2$

There are only two vertices in B and $d_2 = 2$ which means every vertex in a_2 will be adjacent to the vertices in B . This implies $b_1 = 2$, and $b_2 = 0$, and contradicts the assumption $b_2 \neq 0$. So we finished proving the case where $r_b \neq 1$. Therefore we can assume $r_b \geq 2$.

Now assume $a = r_a + 1$:

If $r_a = 1$, then $a = 2$ and it will be the same case as $b = 2$. Assume $r_a \geq 2$, so $a \geq 3$ where $r_a = 2q_b + \epsilon_b + 1 = a - 1$.

$$\begin{aligned}
e &= 2bq_b + 2r_b + b\epsilon_b \\
&= b(2q_b + \epsilon_b) + 2r_b \\
&= b(a - 2) + 2r_b \\
&= ab - 2(b - r_b)
\end{aligned}$$

On the other side,

$$\begin{aligned}
e = 2bq_b + 2r_b + b\epsilon_b &= 2aq_a + 2r_a + a\epsilon_a \\
(b - r_b)(2q_b + \epsilon_b) + r_b(2q_b + \epsilon_b + 2) &= 2aq_a + 2(a - 1) + a\epsilon_a \\
(b - r_b)(a - 2) + ar_b &= a(2q_a + \epsilon_a + 2) - 2 \\
(b - r_b)(a - 2) + ar_b &= a(r_b + 1) - 2 \\
(b - r_b)(a - 2) + ar_b &= ar_b + a - 2 \\
(b - r_b)(a - 2) &= a - 2 \\
(b - r_b - 1)(a - 2) &= 0
\end{aligned}$$

We know $a \geq 3$, $b = r_b + 1$, which means $e = ab - 2(b - r_b) = ab - 2$, which is the bipartite complement of the exception $e = 2$. So we proved $r_b \neq 1$ or $a - r_a \neq 1$. This completes the proof. □

Chapter 4

Results

4.1 Complete Multipartite Graphs

4.1.1 Complete Tripartite Graph

Theorem 4.1. For a complete tripartite graph $K(A, B, C)$ with $|A| = a$, $|B| = b$ and $|C| = c$, assume $a \geq b \geq c$, then the NASCs are:

1. $2 \mid (ab + ac + bc)$
2. $ab \leq ac + bc$

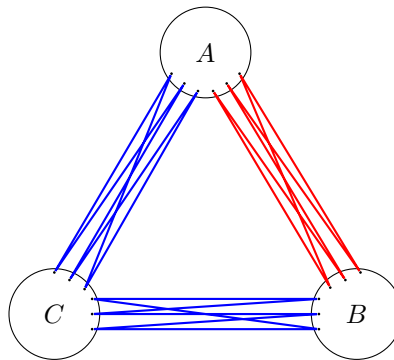


Figure 4.1: $K(A, B, C)$

Necessity:

Proof. $2 \mid (ab + ac + bc)$ comes from the fact that the total number of edges must be divisible by 2 since there are two edges in P_3 . For $ab \leq ac + bc$:

We have three kinds of paths, let x , y , z be the number of the paths $C \rightarrow A \rightarrow B$, $A \rightarrow B \rightarrow C$ and $A \rightarrow C \rightarrow B$ respectively, then

$$x + y = ac$$

$$x + y = ab$$

$$y + z = bc$$

So

$$\begin{aligned} x &= \frac{1}{2}(ac + ab - bc) \Rightarrow bc \leq ac + ab \\ y &= \frac{1}{2}(ab + bc - ac) \Rightarrow ac \leq ab + bc \\ z &= \frac{1}{2}(ac + bc - ab) \Rightarrow ab \leq ac + cb \end{aligned}$$

So if we have $ab \leq ac + cb$, the other two follow easily since $a \geq b \geq c$. □

Sufficiency:

Proof. Let A, B, C be sets of size a, b, c respectively. Assume the necessary conditions are satisfied, then we can find proper x, y, z . Find subgraphs S_1 of $K(C, A)$ and S_2 of $K(A, B)$ with x edges, as in Lemma 3.4, so that their degrees agree on A (thus $S_1 \cup S_2$ is a union of x gregarious paths). Do the same for y paths in $K(A, B) \cup K(B, C)$ and z paths in $K(A, C) \cup K(C, B)$. Now we take the union of these three collections of paths, taking care to rename vertices as in Lemma 3.3. Thus the resulting graph will be the required complete tripartite graph. □

4.2 Star Multipartite Graphs

Definition 4.2. A *star* is a tree consisting of one vertex (called the root) adjacent to the all others. So a star multipartite graph $S = (A; B_1, B_2, \dots, B_n)$ has $|A| = a$ non-adjacent

vertices in the root which are adjacent to all the other sets of vertices (B_1, B_2, \dots, B_n) where $|B_i| = b_i$ for any i .

Theorem 4.3. *Let $S = (A; B_1, B_2, \dots, B_n)$ be a star multipartite graph and assume $b_1 \geq b_2 \geq \dots \geq b_n$. The NASCs are:*

1. $2 \mid (b_1 + b_2 + \dots + b_n)$
2. $b_1 \leq b_2 + b_3 + \dots + b_n$

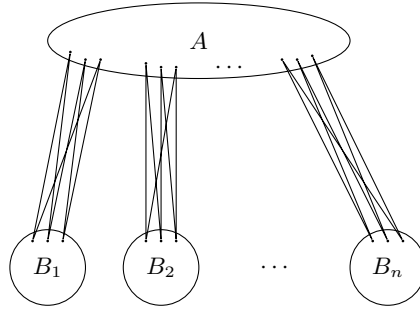


Figure 4.2: Multipartite Star $S(A; B_1, B_2, \dots, B_n)$

Necessity:

Proof. Let v be a vertex in A . So all the gregarious paths passing through v have both ends in $B_1 \cup B_2 \cup \dots \cup B_n$. So $2 \mid (b_1 + b_2 + \dots + b_n)$. For the second condition, the number of the vertices in any b_i should be less than the number of remaining vertices, because if you fix a vertex, say v in a , then the gregarious paths passing through v gives a one-to-one matching in between vertices. So the number of vertices in any part, b_i , should be less than the sum of the number of vertices in the remaining parts. So for any $1 \leq i \leq n$, $b_i \leq b_2 + \dots + b_{i-1} + b_{i+1} + \dots + b_n$. Therefore, if $b_1 \leq b_2 + b_3 + \dots + b_n$ is true, then for any $1 \leq i \leq n$, $b_i \leq b_2 + \dots + b_{i-1} + b_{i+1} + \dots + b_n$ is also true since $b_1 \geq b_2 \geq \dots \geq b_n$. \square

Sufficiency:

Proof. First, take any vertex v in a , then find the gregarious decomposition of $(v; b_1, b_2, \dots, b_n)$. Afterwards, we put the copies of this decomposition on the remaining vertices in A (every vertex in a has the same degree). To find the gregarious decomposition of $(v; b_1, b_2, \dots, b_n)$:

1. take a P_3 between the first (biggest) two parts.
2. reorder $(b_1 - 1, b_2 - 1, \dots, b_n)$ so it is non-increasing.
3. repeat steps 1 and 2 until there are no edges left.

Now we need to prove that in each step the graph we get still satisfies the necessary conditions. The proof of the first condition is easy since we start with an even number of vertices and in each step we just remove two vertices, so in the next step we should still have an even number of vertices.

Now we need prove that in each step we preserve the second condition. We will use induction. Assume that in the k^{th} step we have $(b_1^{(k)}, b_2^{(k)}, \dots, b_n^{(k)})$. For $k = 1$ the second condition holds since $(b_1^{(1)}, b_2^{(1)}, \dots, b_n^{(1)}) = (b_1, b_2, \dots, b_n)$ and for any $1 \leq i \leq n$, $b_i \leq b_2 + \dots + b_{i-1} + b_{i+1} + \dots + b_n$.

To use induction, assume the condition holds for k :

for any $1 \leq i \leq n$, $b_i^{(k)} \leq b_2^{(k)} + \dots + b_{i-1}^{(k)} + b_{i+1}^{(k)} + \dots + b_n^{(k)}$.

So, we need to prove it holds for $k + 1$:

for any $1 \leq i \leq n$, $b_i^{(k+1)} \leq b_2^{(k+1)} + \dots + b_{i-1}^{(k+1)} + b_{i+1}^{(k+1)} + \dots + b_n^{(k+1)}$

Fix i , $1 \leq i \leq n$.

Case 1: $b_i^{(k+1)} = b_i^{(k)} - 1$:

If we remove one vertex from b_i^k , there exists an m with $1 \leq m \leq n$ such that $b_m^{(k+1)} = b_m^k - 1$.

In addition, $b_j^{(k+1)} = b_j^k$ for any j except $j = m, i$. So,

$$\begin{aligned} b_i^{(k+1)} = b_i^{(k)} - 1 &\leq b_1^{(k)} + \dots + b_m^{(k)} - 1 + \dots + b_n^{(k)} \\ &\leq b_1^{(k+1)} + \dots + b_m^{(k+1)} + \dots + b_n^{(k+1)} \end{aligned}$$

Case 2: $b_i^{(k+1)} = b_i^{(k)}$:

Since we removed two vertices in each step, there exist $b_p^{(k)}$ and $b_q^{(k)}$ such that $b_p^{(k)} \geq b_q^{(k)} \geq b_i^{(k)}$.

If $b_i^{(k)} \geq 2$, then $b_i^{(k+1)} = 2 \leq b_p^{(k+1)} + b_q^{(k+1)} + \dots$.

If $b_i^{(k)} = 1$ and $b_p^{(k)} = b_q^{(k)} = 1$, then we should have at least one more $b_w^{(k)} = 1$ since we have an even number of vertices in each step. So

$$\begin{aligned} b_i^{(k+1)} = 1 &\leq b_p^{(k+1)} + b_q^{(k+1)} + b_w^{(k+1)} + \dots \\ &\leq b_p^{(k)} - 1 + b_q^{(k)} - 1 + b_w^{(k)} + \dots \\ &\leq 1 - 1 + 1 - 1 + 1 + \dots \\ &\leq 1 + \dots \end{aligned}$$

□

Note that the following are equivalent:

1. There is a gregarious P_3 decomposition of $S(A; B_1, B_2, \dots, B_n)$.
2. There is a loopless multigraph with degree sequence (b_1, b_2, \dots, b_n) .
3. The complete multigraph $K(B_1, B_2, \dots, B_n)$ has a perfect matching.

4.3 Cycle Multipartite Graphs

4.3.1 Even Cycles

Theorem 4.4. *For an even cycle multipartite graph $C(B_1, \dots, B_{2n})$, the NASCs are:*

1. $b_1b_2 + b_3b_4 + \dots + b_{2n-1}b_{2n} = b_2b_3 + b_4b_5 + \dots + b_{2n}b_1$
2. for any $1 \leq i \leq 2n$, $b_ib_{i+1} \leq b_{i-1}b_i + b_{i+1}b_{i+2}$

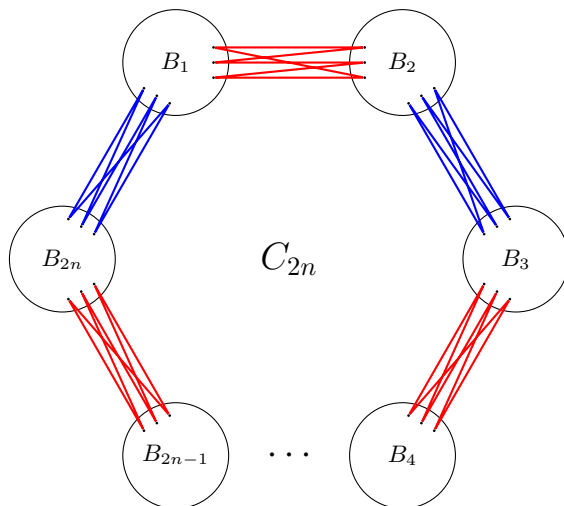


Figure 4.3: Multipartite Even Cycle C_{2n}

Necessity:

Proof. For any $1 \leq i \leq 2n$ let x_i be the number of gregarious paths that have their middle vertex in B_i . Then,

$$\begin{aligned}
 x_1 + x_2 &= b_1 b_2 \\
 x_2 + x_3 &= b_2 b_3 \\
 &\vdots \\
 x_{2n-1} + x_{2n} &= b_{2n-1} b_{2n} \\
 x_{2n} + x_1 &= b_{2n} b_1
 \end{aligned}$$

If we add the first, third, fifth, ..., and $(2n - 1)^{th}$ equations, we get,

$$x_1 + x_2 + \cdots + x_{2n} = b_1 b_2 + b_3 b_4 + \cdots + b_{2n-1} b_{2n}$$

and if we add the second, fourth, sixth, ..., and $(2n)^{th}$ equations and rearrange the x_i 's, we get:

$$x_1 + x_2 + \cdots + x_{2n} = b_2b_3 + b_4b_5 + \cdots + b_{2n}b_1$$

So these two equations give the first condition. For the second condition, let $x_1 = x$ and $x \geq 0$, then:

$$x_2 = b_1b_2 - x$$

$$x_3 = b_2b_3 - x_2$$

$$x_4 = b_3b_4 - x_3$$

$$\vdots$$

$$x_{2n} = b_{2n-1}b_{2n} - x_{2n-1}$$

and if we get all equations in terms of x ,

$$x_2 = b_1b_2 - x$$

$$x_3 = b_2b_3 - b_1b_2 + x$$

$$x_4 = b_3b_4 - b_2b_3 + b_1b_2 - x$$

$$\vdots$$

$$x_{2n} = b_{2n-1}b_{2n} - b_{2n-2}b_{2n-1} + \cdots + x$$

If we use $x \geq 0$ and the equations we have above, then we get:

$$b_i b_{i+1} \leq b_{i-1} b_i + b_{i+1} b_{i+2} \text{ for any } 1 \leq i \leq 2n. \quad \square$$

Sufficiency:

Proof. If the necessary conditions are satisfied, we can find all the x_i 's for $1 \leq i \leq 2n$, then we use the same technique that we used in the proof of Theorem 4.1 to find gregarious a P_3 decomposition. \square

4.3.2 Odd Cycles

Theorem 4.5. *For an odd cycle multipartite graph $C(B_1, \dots, B_{2n+1})$, the NASCs are:*

1. $2 \mid \sum_{i=1}^{i=2n} b_i b_{i+1}$ (# of the edges)

2. for any $1 \leq i \leq 2n + 1$, $b_i b_{i+1} \leq b_{i-1} b_i + b_{i+1} b_{i+2}$

3. for any $1 \leq i \leq 2n + 1$,

$$b_{i+1} b_{i+2} + b_{i+3} b_{i+4} + \dots + b_{i+2n-2} b_{(i+1)+2n-2} \leq b_i b_{i+1} + b_{i+2} b_{i+3} + \dots + b_{i+2n} b_{(i+1)+2n}$$

where the subscripts of the b 's are taken $(\text{mod } 2n + 1)$.

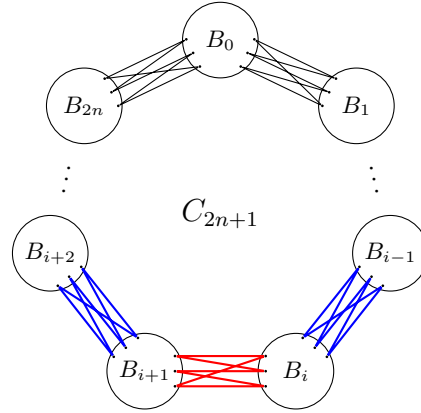


Figure 4.4: Multipartite Odd Cycle C_{2n+1}

Necessity:

Proof. The first condition comes from the fact that the total number of edges is divisible by

2. To get the second condition, for any $1 \leq i \leq 2n + 1$ let x_i be the number of gregarious

paths that have their middle vertex in B_i . Then

$$\begin{aligned}
 x_1 + x_2 &= b_1 b_2 \\
 x_2 + x_3 &= b_2 b_3 \\
 &\vdots \\
 x_{2n} + x_{2n+1} &= b_{2n} b_{2n+1} \\
 x_{2n+1} + x_1 &= b_{2n+1} b_1
 \end{aligned}$$

To find x_1 ;

$$\begin{aligned}
 (+)x_1 + x_2 &= b_1 b_2 \\
 (-)x_2 + x_3 &= b_2 b_3 \\
 &\vdots \\
 (-)x_{2n} + x_{2n+1} &= b_{2n} b_{2n+1} \\
 (+)x_{2n+1} + x_1 &= b_{2n+1} b_1
 \end{aligned}$$

then we get

$$\begin{aligned}
 x_1 &= \frac{b_1 b_2 - b_2 b_3 + b_3 b_4 - \cdots - b_{2n} b_{2n+1} + b_{2n+1} b_1}{2} \\
 &= \frac{b_1 b_2 + b_3 b_4 + \cdots + b_{2n+1} b_1 - (b_2 b_3 + \cdots + b_{2n} b_{2n+1})}{2}
 \end{aligned}$$

Using the same technique we can get all the x_i 's along with condition 3 since each $x_i \geq 0$.

Condition 2 is the same as the even cycle case. \square

Sufficiency:

Proof. After finding x_i , constructing the gregarious P_3 decomposition is the same as for the even cycle case. \square

4.4 Path Multipartite Graphs

Theorem 4.6. For a path multipartite graph $P(B_1, B_2, \dots, B_n)$, the NASCs are:

1. $b_3 \geq b_1$ and $b_{n-2} \geq b_n$
2. $b_1 b_2 + b_3 b_4 \cdots + b_{k-1} b_k = b_2 b_3 + b_4 b_5 + \cdots + b_{l-1} b_l$
3. for any $2 \leq i \leq n-2$, $b_i b_{i+1} \leq b_{i-1} b_i + b_{i+1} b_{i+2}$

where k is the largest **even** number such that $k \leq n$ and l is the largest **odd** number such that $l \leq n$.

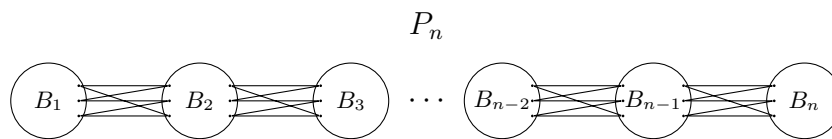


Figure 4.5: Multipartite Path P_n

Necessity:

Proof. For any $2 \leq i \leq n - 1$ let x_i be the number of gregarious paths that have the middle vertex in B_i .

$$\begin{aligned}x_2 &= b_1 b_2 \\x_2 + x_3 &= b_2 b_3 \\x_3 + x_4 &= b_3 b_4 \\&\vdots \\x_{n-2} + x_{n-1} &= b_{n-2} b_{n-1} \\x_{n-1} &= b_{n-1} b_n\end{aligned}$$

The first condition comes from the fact that $x_3 = b_2 b_3 - x_2 = b_2 b_3 - b_1 b_2 = b_2(b_3 - b_1)$. So we get $b_3 \geq b_1$ since $x_2 \geq 0$. We can get $b_n \geq b_{n-2}$ in the same way. We can find the remaining x_i 's easily.

If we add the first, third, fifth,..., and $(k - 1)^{th}$ equations, we get,

$$x_2 + x_3 \cdots + x_{n-1} = b_1 b_2 + b_3 b_4 + \cdots + b_{k-1} b_k$$

and if we add the second, fourth, sixth,..., and $(l - 1)^{th}$ equations, we get:

$$x_2 + x_3 + \cdots + x_{n-1} = b_2 b_3 + b_4 b_5 + \cdots + b_{l-1} b_l$$

So these two equations give the second condition. The third condition is the same as the condition in the cycle case. □

Sufficiency:

Proof. If the necessary conditions are satisfied, we can find all the x_i 's for $2 \leq i \leq n - 1$, then we use the same technique that we used in the proof of Theorem 4.1 to find gregarious a P_3 decomposition. \square

4.5 Some Tree Multipartite Graphs

Definition 4.7. Let $T(C_1, \dots, C_m; A_1, A_2; B_1, \dots, B_n)$ be a multipartite graph such that two multipartite stars $S(A_1; C_1, \dots, C_m)$ and $S(A_2; B_1, \dots, B_n)$ are attached to each other via putting a complete bipartite graph on bipartition (A_1, A_2) . See figure 4.6.

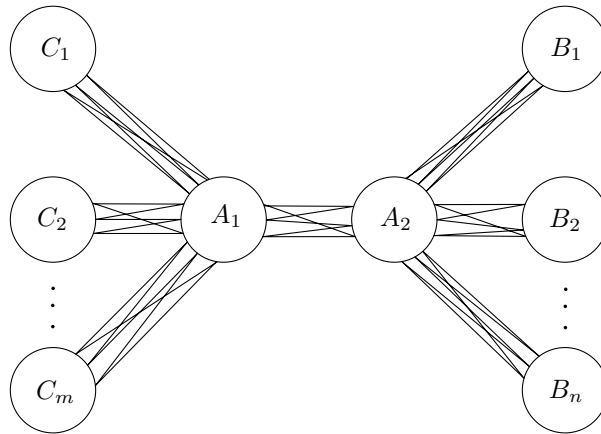


Figure 4.6: $T(C_1, \dots, C_m; A_1, A_2; B_1, \dots, B_n)$

Definition 4.8. Define $T(A_1, A_2, A_3; B_1, \dots, B_n)$ by using definition 4.7 as $T(A_1; A_2, A_3; B_1, \dots, B_n)$. See figure 4.7.

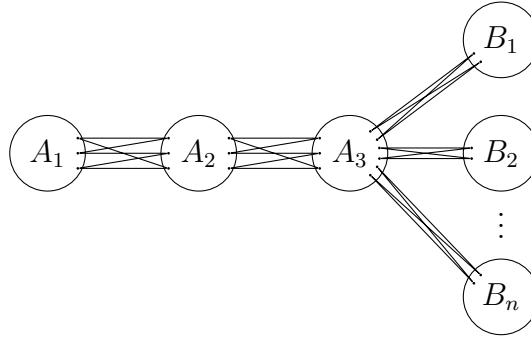


Figure 4.7: $T(A_1, A_2, A_3; B_1, \dots, B_n)$

Lemma 4.9. Let $G = (E, V)$ be a graph. There is an orientation of G such that for all $v \in V$, $|out(v) - in(v)| \leq 1$.

Proof. We can assume that G is connected.

Case 1: If all vertices have even degree, then there exists an Euler trail, we can orient the graph this way.

Case 2: If G has some vertices with odd degree, make an extra vertex u and connect all those vertices to u , then find an Euler trail on $G \cup \{u\}$ and remove the edges at the end. For all $v \in V$, we still have $|out(v) - in(v)| \leq 1$ since we remove one edge from each vertex with odd degree .

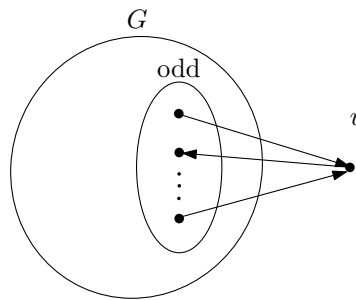


Figure 4.8: Orientation of G

□

Theorem 4.10. [14] *Let A, B, I be finite non-empty sets, let $f : B \times I \rightarrow \mathbb{N}$ be such that for all $t \in B$, $\sum_{i \in I} f(t, i) = |A|$. Then the edges of $K(A, B)$ can be partitioned into spanning subgraphs G_i , $i \in I$, such that for each $i \in I$, G_i is balanced on A , and for each $t \in B$, the degree of t in G_i is $f(t, i)$.*

Proof. If $|A| = 1$, then the proof is trivial.

Now suppose $|A| = 2$. Let $A = \{s_1, s_2\}$. Form a graph H on vertex set I as follows:

For each $t \in B$, H has an edge e_t : If $f(t, i) = 2$, (and so $f(t, j) = 0$ for all other $j \in I$) then e_t is a loop at vertex i of H . If $f(t, i) = 1 = f(t, j)$, $i \neq j$, ($f(t, k) = 0$ for the other $k \in I$), then e_t joins the vertices i and j in H .

Orient H so that at each vertex of H the indegree and outdegree differ by at most 1 using Lemma 4.9. For each $t \in B$, if e_t is directed from i to j in the oriented H , place the edge between t and s_1 in G_i , and the edge between t and s_2 in G_j (see example 4.11).

If $|A| \geq 3$, then partition the edges of $K(A, B)$ into spanning subgraphs G_i whose degrees on B are given by f (this is certainly possible by the sum condition on f). If everything is balanced on $|A|$, then we are done. Otherwise degrees in some G_i differ by 2 or more. Fix i , and let s_1, s_2 be two vertices in A , whose degrees differ by 2 or more in G_i . So use the previous case where $|A| = 2$ on this graph to find the balanced distribution. Using this method repeatedly for each unbalanced pair of vertices of G_i in A , finally we can get the balanced distribution on A . Afterwards, we can repeat the same process for the other G_j for each $j \in I$.

Now we need to prove that this process will stop after finitely many steps. Let $V_1 = (a_1, \dots, a_i, \dots, a_j, \dots, a_n)$ be a integer vector with fixed sum $\sum a_i = a$. So the shortest integer vector with respect to the Euclidean metric with the fixed sum of the entries is the balanced one. To see this assume $a_j \geq a_i + 2$, then if we balance a_i and a_j , we get

$V_2 = (a_1, \dots, a_i + 1, \dots, a_j - 1, \dots, a_n)$ and $|V_2|^2 \leq |V_1|^2 - 2$ since,

$$\begin{aligned}
|V_2|^2 &= a_1^2 + \dots + (a_i + 1)^2 + \dots + (a_j - 1)^2 + \dots + a_n^2 \\
&= a_1^2 + \dots + a_i^2 + 2a_i + 1 + \dots + a_j^2 - 2a_j + 1 + \dots + a_n^2 \\
&= (a_1^2 + \dots + a_i^2 + \dots + a_j^2 + \dots + a_n^2) + 2(a_i - a_j) + 2 \\
&= |V_1|^2 + 2(a_i - a_j) + 2 \\
&\leq |V_1|^2 + 2(-2) + 2 \\
&\leq |V_1|^2 - 2
\end{aligned}$$

This means that when we balance a pair of entries in the vector at a time, the vector gets shorter, and after finitely many steps we will find the shortest one. This completes the proof. □

Example 4.11. Let G be a bipartite graph on bipartition (A, B) where $A = \{s_1, s_2\}$ and $B = \{v_1, v_2, v_3, v_4\}$. Let $I = \{\text{green, blue, red}\}$. We want to get **green** and **blue** balanced on A without changing the color census on B (see the first picture in Figure 4.9). Then using the method defined in Theorem 4.10 build a graph H (see the second picture in Figure 4.9) and orient H so that $|in(w) - out(w)| \leq 1$ for every vertex w in H . Finally, we can swap edges of G with respect to the orientation on H to get a balanced coloring on A .

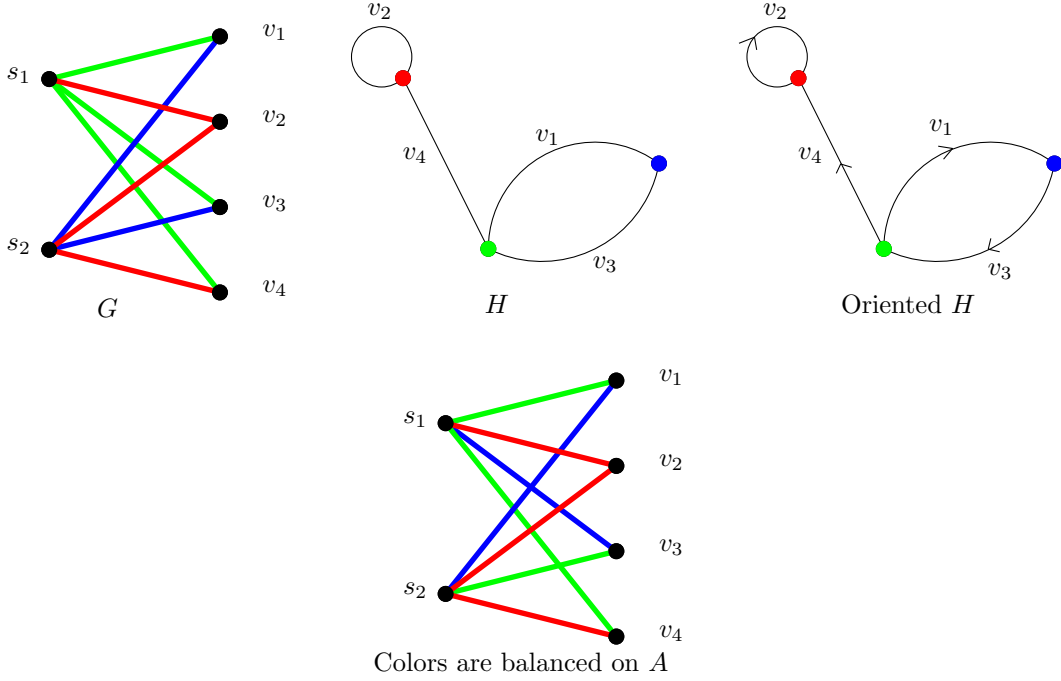


Figure 4.9: An Example for Theorem 4.10

4.5.1 Necessary And Sufficient Conditions for $T(A_1, A_2, A_3; B_1, \dots, B_n)$

Theorem 4.12. For a graph $T(A_1, A_2, A_3; B_1, \dots, B_n)$ assume $b_n \leq \dots \leq b_2 \leq b_1$, the NASCs are:

1. $2 \mid [a_2(a_1 + a_3) + a_3(b_1 + b_2 + \dots + b_n)]$
2. $a_1 \leq a_3$
3. $a_1 a_2 + b_1 a_3 \leq a_3(a_2 + b_2 + b_3 + \dots + b_n)$
4. $a_2 a_3 \leq a_1 a_2 + a_3(b_1 + b_2 + \dots + b_n)$
5. If
 - $a_2 + (b_1 + \dots + b_n)$ is even, then $a_1 a_2$ is even.
 - $a_2 + (b_1 + \dots + b_n)$ is odd, then $a_1 a_2 - a_3$ is even and non-negative.

Necessity:

Proof. Let $G = T(A_1, A_2, A_3; B_1, \dots, B_n)$. Condition 1 comes from the fact that the number of edges is even. We can get conditions 2 - 4 using Tutte's f-factor Theorem on $L(G)$. $L(G)$ is union of a complete graph on $n + 1$ vertices and an edge attached at the vertex A_2A_3 . If we check all the possible $L(G)$ triples $B = (S, T, U)$ where T is independent and $\lambda(T, U) = 0$ from theorem 2.12, all will reduce to the the following three cases in figure 4.10. We need to check $k(B, h) + h(T) \leq h(S)$ for each case. From the first picture in the figure 4.10, we will get $a_2a_3 \leq a_1a_2$ which gives condition 2. From the second picture, we get $a_1a_2 + b_1a_3 \leq a_2a_3 + b_2a_3 + \dots + b_na_3$ which gives condition 3. From the last picture, we get $a_2a_3 \leq a_1a_2 + a_3b_1 + a_3b_2 + \dots + a_3b_n$ which gives condition 4.

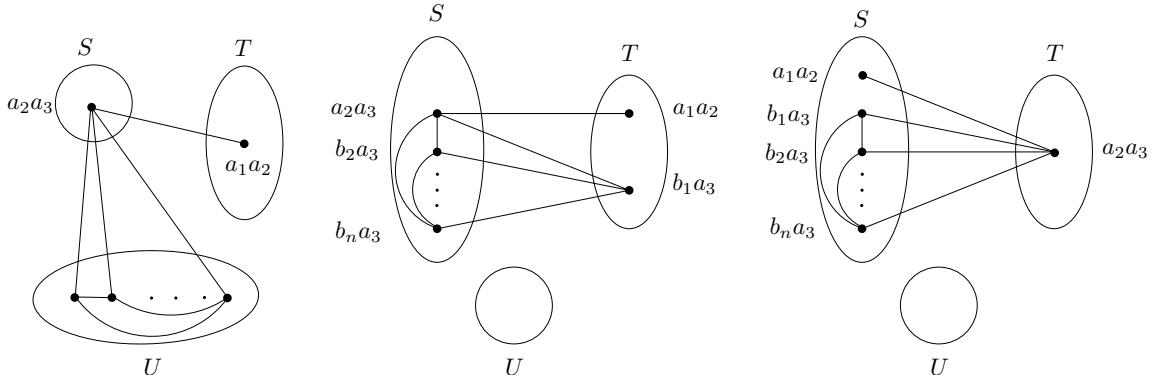


Figure 4.10: How to get conditions 2 - 4

For condition 5, we need to consider all the types of paths we have and the degree of any vertex in A_3 . Firstly, the degree of any vertex v in A_3 is $deg(v) = a_2 + b_1 + \dots + b_n$. Let x_1 be the number of paths passing through the sets of vertices $A_1 \rightarrow A_2 \rightarrow A_3$ so $x_1 = a_1a_2$. In the same way, $y_i : A_2 \rightarrow A_3 \rightarrow B_i$ for any $1 \leq i \leq n$, and $w_{ij} : B_i \rightarrow A_3 \rightarrow B_j$ for any $1 \leq i \leq j \leq n$. Here, both y_i and w_{ij} have their middle vertices in A_3 , so if $deg(v)$ is even, then x_1 must be even. If $deg(v)$ is odd, then we should have enough x_1 type paths which means $a_1a_2 \geq a_3$. That gives $a_1a_2 - a_3$ non-negative. To see that $a_1a_2 - a_3$ is even, consider the vertices in A_3 and distribution of a_1a_2 edges on A_3 . There are α_i 's for

$1 \leq i \leq a_3$ such that $\sum_{i=1}^{a_3} \alpha_i = a_1 a_2$ where each $\alpha_i = 2\beta_i + 1$, an odd number. $\beta_i = \frac{\alpha_i - 1}{2}$.
 $\sum_{i=1}^{a_3} \beta_i = \sum_{i=1}^{a_3} \frac{\alpha_i - 1}{2} = \frac{1}{2}(a_1 a_2 - a_3)$. Therefore $a_1 a_2 - a_3$ is an even number. \square

Sufficiency:

Proof. If the necessary conditions are satisfied we can find proper x_1 , y_i 's and w_{ij} 's. In between pairs of sets (A_1, A_2) and (A_3, B_i) for any i , we can find balanced edge distributions with the required numbers as we did for the sufficiency case of the Theorem 4.1. The only problem is finding a construction for (A_2, A_3) since we need to find a parity balanced distribution of x_1 edges on A_3 with respect to the parity of $a_2 + (b_1 + \dots + b_n)$ (see condition 5 in Theorem 4.12). We also need to find a balanced distribution for the remaining y_i 's. To be able to find this special distribution we can use theorem 4.10 and choose the degrees on A_3 to get balanced degrees on A_2 . \square

4.5.2 Necessary And Sufficient Conditions for $T(C_1, \dots, C_m; A_1, A_2; B_1, \dots, B_n)$

Theorem 4.13. For a graph $T(C_1, \dots, C_m; A_1, A_2; B_1, \dots, B_n)$ assume $c_m \leq \dots \leq c_2 \leq c_1$ and $b_n \leq \dots \leq b_2 \leq b_1$, and let $C = C_1 \cup C_2 \cup \dots \cup C_m$ and $|C| = c$, $B = B_1 \cup B_2 \cup \dots \cup B_n$, $|B| = b$ and $d_1 = c + a_2$, $d_2 = b + a_1$. The NASCs are:

1. $2 \mid [a_1(c_1 + c_2 + \dots + c_m) + a_1 a_2 + a_2(b_1 + b_2 + \dots + b_n)]$

2. $c_1 \leq a_2 + (c_2 + \dots + c_m)$ and $b_1 \leq a_1 + (b_2 + \dots + b_n)$

3. $a_1 c_1 + a_2 b_1 \leq a_1(c_2 + \dots + c_m) + a_1 a_2 + a_2(b_2 + \dots + b_n)$

4. $a_1 a_2 \leq a_1(c_1 + c_2 + \dots + c_m) + a_2(b_1 + b_2 + \dots + b_n)$

5. If

- d_1 and d_2 are even, then $a_1 a_2$ is even.

- d_1 is even and d_2 is odd, then either both a_1 and a_2 are odd, both even or a_1 is odd and a_2 even. In addition, $ca_1 - a_2 \geq 0$.
- d_1 is odd and d_2 is even, then either both a_1 and a_2 are odd, both even or a_1 is even and a_2 odd. In addition, $ba_2 - a_1 \geq 0$.
- d_1 and d_2 are odd, then both a_1 and a_2 are even. In addition, $ca_1 - a_2 \geq 0$ and $ba_2 - a_1 \geq 0$.

with the following exceptions:

d_1	d_2	a_1	a_2	c_1 & b_1
even	even	2	2	any c_1 with $1 + (b_2 + \dots + b_n) \leq b_1 \leq 2 + (b_2 + \dots + b_n)$
		even	1	any c_1 with $1 + (b_2 + \dots + b_n) \leq b_1 \leq a_1 + (b_2 + \dots + b_n)$
		1	even	any b_1 with $1 + (c_2 + \dots + c_m) \leq c_1 \leq a_2 + (c_2 + \dots + c_m)$
even	odd	odd	1	any c_1 with $1 + (b_2 + \dots + b_n) \leq b_1 \leq a_1 + (b_2 + \dots + b_n)$
		1	even	any c_1 with $b_1 = 1 + (b_2 + \dots + b_n)$
odd	even	1	odd	any b_1 with $1 + (c_2 + \dots + c_m) \leq c_1 \leq a_2 + (c_2 + \dots + c_m)$
		even	1	any b_1 with $c_1 = 1 + (c_2 + \dots + c_m)$

Table 4.1: Exceptions for Theorem 4.13

Necessity:

Proof. Let $G = T(C_1, \dots, C_m; A_1, A_2; B_1, \dots, B_n)$. Condition 1 comes from the fact that the number of edges is even. We can get conditions 2 - 4 using Tutte's f-factor Theorem on $L(G)$. $L(G)$ is union of two complete graphs on m and n vertices attached at the vertex A_1A_2 . If we check all the possible $L(G)$ triples $B = (S, T, U)$ where T is independent and $\lambda(T, U) = 0$ from theorem 2.12, all will reduce to the the following three cases in figure 4.11. We need to check $k(B, h) + h(T) \leq h(S)$ for each case. From the first picture in the figure 4.11, we will get $c_1 \leq a_2 + (c_2 + \dots + c_m)$ and in the same picture if we replace B 's with C 's

and C 's with B 's then we get $b_1 \leq a_1 + (b_2 + \dots + b_n)$ which gives condition 2. From the second picture, we get $a_1c_1 + a_2b_1 \leq a_1(c_2 + \dots + c_m) + a_1a_2 + a_2(b_2 + \dots + b_n)$ which gives condition 3. From the last picture, we get $a_1a_2 \leq a_1(c_1 + c_2 + \dots + c_m) + a_2(b_1 + b_2 + \dots + b_n)$ which gives condition 4.

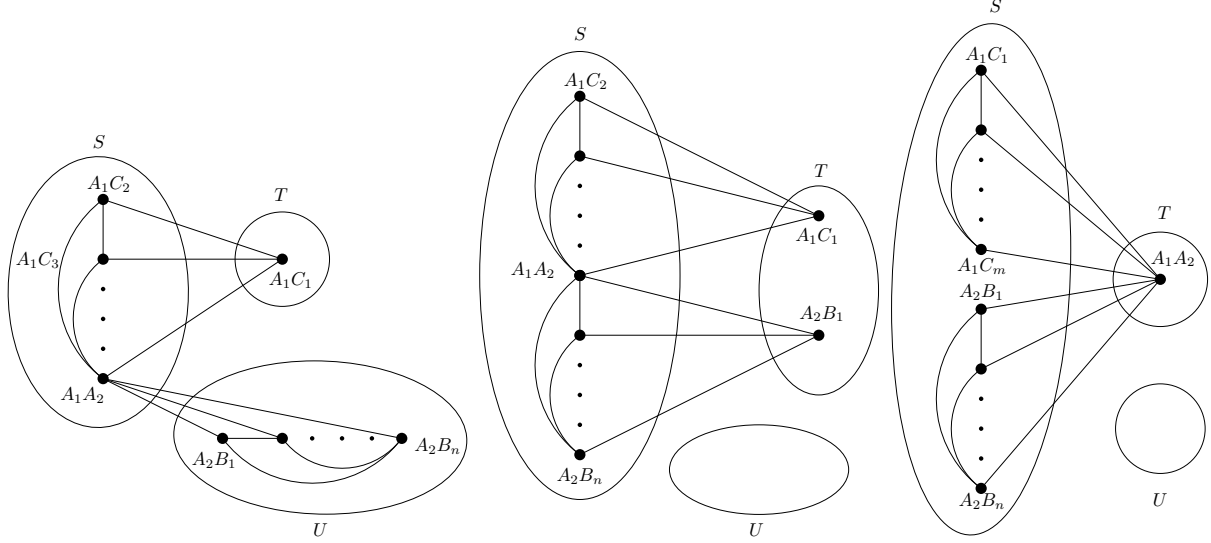


Figure 4.11: How to get conditions 2 - 4

For condition 5, we need to consider all the types of paths we have and the degree of any vertex in A_1 and A_2 . Firstly, the degree of any vertex v_1 in A_1 is:

$$d_1 = \deg(v_1) = a_2 + c_1 + \dots + c_m = a_2 + c$$

and the degree of any vertex v_2 in A_2 is:

$$d_2 = \deg(v_2) = a_1 + b_1 + \dots + b_n = a_1 + b.$$

Let x_i be the number of paths passing through the sets of vertices $C_i \rightarrow A_1 \rightarrow A_2$ for any $1 \leq i \leq m$ and let $x = \sum_{i=1}^m x_i$. In the same way, $y_j : B_j \rightarrow A_2 \rightarrow A_1$ for any $1 \leq j \leq n$ and $y = \sum_{j=1}^n y_j$. So $x + y = a_1a_2$. In addition, we have $w_{ij} : C_i \rightarrow A_1 \rightarrow C_j$ for any $1 \leq i \leq j \leq m$ and $z_{kl} : B_k \rightarrow A_2 \rightarrow B_l$ for any $1 \leq k \leq l \leq n$. Here w_{ij} 's have their middle vertex in A_1 and z_{kl} 's have their middle vertex in A_2 , so we have four cases with respect to the parity of d_1 and d_2 . So the parity of x and d_2 , and y and x_1 must be consistent (see figure 4.12).

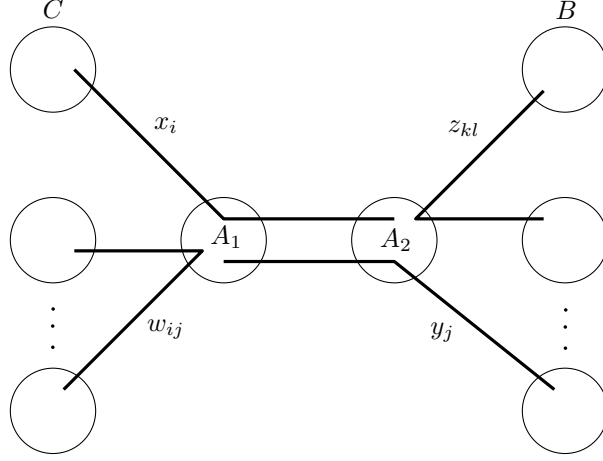


Figure 4.12: Types of paths

Case 1: If d_1 and d_2 are even, then y and x are even so a_1a_2 is even since $x + y = a_1a_2$.

Case 2: If d_1 is even and d_2 is odd, then y is even and $x \equiv a_2 \pmod{2}$ and $x \geq a_2$. So we can get $ca_1 - a_2 \geq 0$ since $ca_1 \geq x$. To get either both a_1 and a_2 odd, both even or a_1 is odd and a_2 even, see:

$$\begin{aligned} x + y &= a_1a_2 \\ x + y &\equiv a_1a_2 \pmod{2} \\ x &\equiv a_1a_2 \pmod{2} \text{ since } y \text{ is even} \end{aligned}$$

Case 3: If d_1 is odd and d_2 is even, then this is the same as case 2, just replace a_2 with a_1 .

Case 4: If d_1 is odd and d_2 is odd, then $y \equiv a_1 \pmod{2}$ and $y \geq a_1$, and $x \equiv a_2 \pmod{2}$ and $x \geq a_2$. We can get $ca_1 - a_2 \geq 0$ and $ba_2 - a_1 \geq 0$ in the same way as in case 2. To get both a_1 and a_2 even, see:

$$\begin{aligned} x + y &= a_1a_2 \\ x + y &\equiv a_1a_2 \pmod{2} \\ a_1 + a_2 &\equiv a_1a_2 \pmod{2} \text{ since } x \equiv a_2 \pmod{2} \text{ and } y \equiv a_1 \pmod{2} \end{aligned}$$

$a_1 + a_2 \equiv a_1 a_2 \pmod{2}$ is only satisfied when both a_1 and a_2 are even.

Note that theorem 4.12 is a special case of theorem 4.13. In theorem 4.13, if we get $c_2 = c_3 = \dots = c_m = 0$ and replace c_1 with a_1 , a_1 with a_2 and a_2 with a_3 we will get exactly the same conditions as in theorem 4.12. \square

Sufficiency:

Proof. If the necessary conditions are satisfied we can find proper x_i 's, y_j 's, w_{ij} 's and v_{kl} 's. First we find a proper x and y then we will find w_{ij} 's and v_{kl} 's since we have more restriction on x and y . In between pairs of sets (C_i, A_1) for any $1 \leq i \leq m$, and (A_2, B_j) for any $1 \leq j \leq n$, we can find balanced edge distributions with the required numbers as we did for the sufficiency case of theorem 4.1. The only problem is finding a construction for (A_1, A_2) since we need to find a parity balanced distribution of $x + y$ edges on A_1 and A_2 with respect to the parity of d_1 and d_2 (see condition 5 in theorem 4.13). To find this parity balanced distribution, we will use theorem 3.12.

Case 1: Assume d_1 and d_2 are even, then y and x are even so $a_1 a_2$ is even since $x + y = a_1 a_2$. So there are three cases for (a_1, a_2) : (even, even), (even, odd) and (odd, even).

If a_1 and a_2 are both even, then we need to find a parity balanced bipartite graph (PBBG) with parameters $(a = a_1, b = a_2, e = x, \epsilon_a = 0, \epsilon_b = 0)$ with bipartite complement $(a = a_1, b = a_2, e = ab - x = y, \epsilon'_a = 0, \epsilon'_b = 0)$ so that the distribution of x on A_2 has even parity ($\epsilon_b = 0$) and the distribution of y on A_1 has even parity ($\epsilon'_a = 0$). We can we can find such a PBBG since the necessary conditions of theorem 3.12 are satisfied.

$$\epsilon_a + \epsilon'_a = 0 + 0 = 0 \equiv b \pmod{2} \text{ and } \epsilon_b + \epsilon'_b = 0 + 0 = 0 \equiv a \pmod{2}$$

$$\epsilon_a a \leq e \leq ab - \epsilon'_a a \Rightarrow 0 \leq x \leq a_1 a_2$$

$$\epsilon_b b \leq e \leq ab - \epsilon'_b b \Rightarrow 0 \leq x \leq a_1 a_2$$

$$\epsilon_a a \equiv ab - \epsilon'_a a \equiv e = x \equiv \epsilon_b b \equiv ab - \epsilon'_b b \equiv 0 \pmod{2}$$

For the exceptions in table 4.1 the only problem concerning this case is when $a = a_1 = \text{even}$, $b = a_2 = \text{even}$, $e = x = 2$, $\epsilon_a = 0$, $\epsilon'_a = 0$, $\epsilon_b = 0$, $\epsilon'_b = 0$. We can solve this problem by choosing $x \geq 4$ since $a_1 \geq 2$ and $a_2 \geq 2$ in the exceptions. In the case of $a_1 = 2 = a_2$, we will not have any y 's which means we need to put a gregarious P_3 decomposition of star multipartite graph on $S = (A_2; B_2, B_2, \dots, B_n)$. The first condition of theorem 4.3 is satisfied since $2 \mid b = b_1 + b_2 + \dots + b_n$ ($d_1 = \text{even} = a_1 + b$ and a_1 is even so b is even). For the second condition of theorem 4.3, we can get $b_1 \leq 2 + b_2 + \dots + b_n$ from condition 2 of theorem 4.13. From here we get two exceptions: $b_1 = 1 + b_2 + \dots + b_n$ and $b_1 = 2 + b_2 + \dots + b_n$. So $T(C_1, \dots, C_m; 2, 2; B_1, \dots, B_n)$ doesn't exist when either $b_1 = 1 + b_2 + \dots + b_n$ or $b_1 = 2 + b_2 + \dots + b_n$.

If a_1 is even and a_2 is odd, then we need to find a PBBG with parameters $(a = a_1, b = a_2, e = x, \epsilon_a = 1, \epsilon_b = 0)$ with bipartite complement $(a = a_1, b = a_2, e = ab - x = y, \epsilon'_a = 0, \epsilon'_b = 0)$ so that the distribution of x edges on a_2 has even parity ($\epsilon_b = 0$) and the distribution of y edges on a_1 has even parity ($\epsilon'_a = 0$). We can find such a PBBG since the necessary conditions of theorem 3.12 are satisfied.

$$\begin{aligned} \epsilon_a + \epsilon'_a &= 1 + 0 = 1 \equiv b \pmod{2} \text{ and } \epsilon_b + \epsilon'_b = 0 + 0 = 0 \equiv a \pmod{2} \\ \epsilon_a a &\leq e \leq ab - \epsilon'_a a \Rightarrow a_1 \leq x \leq a_1 a_2 \\ \epsilon_b b &\leq e \leq ab - \epsilon'_b b \Rightarrow 0 \leq x \leq a_1 a_2 \\ \epsilon_a a &\equiv ab - \epsilon'_a a \equiv e = x \equiv \epsilon_b b \equiv ab - \epsilon'_b b \equiv 0 \pmod{2} \end{aligned}$$

If we check the exceptions in table 4.1, then we see $(a = \text{even} \geq 4, b = \text{odd} \geq 3, e = 2, \epsilon_a = 1, \epsilon'_a = 0, \epsilon_b = 0, \epsilon'_b = 0)$. However, in this case $e = x \geq a_1$ and $a_1 \geq 4$ so we don't have any exception for $e = 2$. There are other exceptions coming from $x \geq a_1$. If $a_2 = 1$, then we don't have any y 's which means we need to put a gregarious P_3 decomposition of a star multipartite graph on $S = (A_2 = 1; B_2, B_2, \dots, B_n)$. The first condition of theorem 4.3 is satisfied since $2 \mid b = b_1 + b_2 + \dots + b_n$ ($d_2 = \text{even} = a_1 + b$ and a_1 is even so b is even). For the second

condition of theorem 4.3, we can get $b_1 \leq a_1 + (b_2 + \dots + b_n)$ from condition 2 of theorem 4.13. From here we get exceptions when $1 + (b_2 + \dots + b_n) \leq b_1 \leq a_1 + (b_2 + \dots + b_n)$. So $T(C_1, \dots, C_m; A_1, 1; B_1, \dots, B_n)$ doesn't exist when $1 + (b_2 + \dots + b_n) \leq b_1 \leq a_1 + (b_2 + \dots + b_n)$.

If a_1 is odd and a_2 is even, then this case is the same as the previous case, just switch a_2 with a_1 .

Case 2: If d_1 is even and d_2 is odd, then y is even and $x \equiv a_2 \pmod{2}$ and $x \geq a_2$. So there are three cases for (a_1, a_2) : (even, even), (odd, odd) and (odd, even).

If a_1 and a_2 are both even, then x and y are both even. We need to find a PBBG with parameters $(a = a_1, b = a_2, e = x, \epsilon_a = 0, \epsilon_b = 1)$ with bipartite complement $(a = a_1, b = a_2, e = a_1a_2 - x = y, \epsilon'_a = 0, \epsilon'_b = 1)$ so that the distribution of x on A_2 has odd parity ($\epsilon_b = 1$) and the distribution of y on A_1 has even parity ($\epsilon'_a = 0$). We can we can find such a PBBG since the necessary conditions of theorem 3.12 are satisfied.

$$\begin{aligned} \epsilon_a + \epsilon'_a &= 0 + 0 = 0 \equiv b \pmod{2} \text{ and } \epsilon_b + \epsilon'_b = 1 + 1 = 0 \equiv a \pmod{2} \\ \epsilon_a a &\leq e \leq ab - \epsilon'_a a \Rightarrow 0 \leq x \leq a_1 a_2 \\ \epsilon_b b &\leq e \leq ab - \epsilon'_b b \Rightarrow a_2 \leq x \leq a_1 a_2 - a_2 \\ \epsilon_a a &\equiv ab - \epsilon'_a a \equiv e = x \equiv \epsilon_b b \equiv ab - \epsilon'_b b \equiv 0 \pmod{2} \end{aligned}$$

If we check the exceptions in table 4.1, we see that we don't have any exception for this case.

If a_1 and a_2 are both odd, then x is odd and y is even. We need to find a PBBG with parameters $(a = a_1, b = a_2, e = x, \epsilon_a = 1, \epsilon_b = 1)$ with bipartite complement $(a = a_1, b = a_2, e = a_1a_2 - x = y, \epsilon'_a = 0, \epsilon'_b = 0)$ so that the distribution of x on A_2 has odd parity ($\epsilon_b = 1$) and the distribution of y on A_1 has even parity ($\epsilon'_a = 0$). We can we can find such

a PBBG since the necessary conditions of theorem 3.12 are satisfied.

$$\begin{aligned}
\epsilon_a + \epsilon'_a &= 0 + 1 = 1 \equiv b \pmod{2} \text{ and } \epsilon_b + \epsilon'_b = 1 + 0 = 1 \equiv a \pmod{2} \\
\epsilon_a a &\leq e \leq ab - \epsilon'_a a \Rightarrow a_1 \leq x \leq a_1 a_2 \\
\epsilon_b b &\leq e \leq ab - \epsilon'_b b \Rightarrow a_2 \leq x \leq a_1 a_2 \\
\epsilon_a a &\equiv ab - \epsilon'_a a \equiv e = x \equiv \epsilon_b b \equiv ab - \epsilon'_b b \equiv 1 \pmod{2}
\end{aligned}$$

If we check the exceptions in table 4.1, then we see ($a = \text{odd} \geq 3, b = \text{odd} \geq 3, e = 2, \epsilon_a = 1, \epsilon'_a = 0, \epsilon_b = 1, \epsilon'_b = 0$). However, in this case $e = x \geq \max\{a_1, a_2\}$ and $a_1, a_2 \geq 3$ so we don't have any exception for $e = 2$. There are other exceptions coming from $x \geq \max\{a_1, a_2\}$. If $a_2 = 1$, then $x = a_1 = a_1 a_2$ so we don't have any y 's which means we need to put a gregarious P_3 decomposition of a star multipartite graph on $S = (A_2 = 1; B_2, B_2, \dots, B_n)$. The first condition of theorem 4.3 is satisfied since $2 \mid b = b_1 + b_2 + \dots + b_n$ ($d_2 = \text{odd} = a_1 + b$ and a_1 is odd so b is even). For the second condition of theorem 4.3, we can get $b_1 \leq a_1 + (b_2 + \dots + b_n)$ from condition 2 of theorem 4.13. From here we get exceptions when $1 + (b_2 + \dots + b_n) \leq b_1 \leq a_1 + (b_2 + \dots + b_n)$. So $T(C_1, \dots, C_m; A_1, 1; B_1, \dots, B_n)$ doesn't exist when $1 + (b_2 + \dots + b_n) \leq b_1 \leq a_1 + (b_2 + \dots + b_n)$.

If a_1 is odd and a_2 is even, then x and y are both even. We need to find a PBBG with parameters ($a = a_1, b = a_2, e = x, \epsilon_a = 0, \epsilon_b = 1$) with bipartite complement ($a = a_1, b = a_2, e = a_1 a_2 - x = y, \epsilon'_a = 0, \epsilon'_b = 0$) so that the distribution of x on A_2 has odd parity ($\epsilon_b = 1$) and the distribution of y on A_1 has even parity ($\epsilon'_a = 0$). We can we can find such a PBBG since the necessary conditions of theorem 3.12 are satisfied.

$$\begin{aligned}
\epsilon_a + \epsilon'_a &= 0 + 0 = 0 \equiv b \pmod{2} \text{ and } \epsilon_b + \epsilon'_b = 1 + 0 = 1 \equiv a \pmod{2} \\
\epsilon_a a &\leq e \leq ab - \epsilon'_a a \Rightarrow 0 \leq x \leq a_1 a_2 \\
\epsilon_b b &\leq e \leq ab - \epsilon'_b b \Rightarrow a_2 \leq x \leq a_1 a_2 \\
\epsilon_a a &\equiv ab - \epsilon'_a a \equiv e = x \equiv \epsilon_b b \equiv ab - \epsilon'_b b \equiv 0 \pmod{2}
\end{aligned}$$

If we check the exceptions in table 4.1, then we see ($a = \text{odd} \geq 3, b = \text{even} \geq 3, e = 2, \epsilon_a = 0, \epsilon'_a = 0, \epsilon_b = 1, \epsilon'_b = 0$). However, in this case $e = x \geq a_2$ and $a_2 \geq 4$ so we don't have any exception for $e = 2$. There is an exception coming from $x \geq a_2$. If $a_1 = 1$, then $x = a_2 = a_1 a_2$ so we don't have any y 's which means we need to put a gregarious P_3 decomposition of a star multipartite graph on $S = (A_2; B_2, B_2, \dots, B_n)$. The first condition of theorem 4.3 is satisfied since $2 \mid b = b_1 + b_2 + \dots + b_n$ ($d_2 = \text{odd} = a_1 + b$ and a_1 is odd so b is even). For the second condition of theorem 4.3, we can get $b_1 \leq a_1 + (b_2 + \dots + b_n)$ from condition 2 of theorem 4.13. From here we get exceptions when $b_1 = 1 + (b_2 + \dots + b_n)$. So $T(C_1, \dots, C_m; 1, A_2; B_1, \dots, B_n)$ doesn't exist when $b_1 = 1 + (b_2 + \dots + b_n)$.

Case 3: If d_1 is odd and d_2 is even, then this case is the same as case 2, just switch a_1 and a_2 .

Case 4: If d_1 and d_2 are both odd, then x, y, a_1, a_2 are even and $y \geq a_1, x \geq a_2$. We need to find a PBBG with parameters ($a = a_1, b = a_2, e = x, \epsilon_a = 1, \epsilon_b = 1$) with bipartite complement ($a = a_1, b = a_2, e = a_1 a_2 - x = y, \epsilon'_a = 1, \epsilon'_b = 1$) so that the distribution of x on A_2 has odd parity ($\epsilon_b = 1$) and the distribution of y on A_1 has odd parity too ($\epsilon'_a = 1$). We can we can find such a PBBG since the necessary conditions of theorem 3.12 are satisfied.

$$\begin{aligned} \epsilon_a + \epsilon'_a &= 1 + 1 = 0 \equiv b \pmod{2} \text{ and } \epsilon_b + \epsilon'_b = 1 + 1 = 0 \equiv a \pmod{2} \\ \epsilon_a a &\leq e \leq ab - \epsilon'_a a \Rightarrow a_1 \leq x \leq a_1 a_2 - a_1 \\ \epsilon_b b &\leq e \leq ab - \epsilon'_b b \Rightarrow a_2 \leq x \leq a_1 a_2 - a_2 \\ \epsilon_a a &\equiv ab - \epsilon'_a a \equiv e = x \equiv \epsilon_b b \equiv ab - \epsilon'_b b \equiv 0 \pmod{2} \end{aligned}$$

If we check the exceptions in table 4.1, we see that we don't have any exception for this case. □

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